# Rectilinear approximation and volume estimates for hereditary bodies via [0,1]-decorated containers 

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#### Abstract

We use the hypergraph container theory of Balogh-Morris-Samotij and Saxton-Thomason to obtain general rectilinear approximations and volume estimates for sequences of bodies closed under certain families of projections. We give a number of applications of our results, including a multicolour generalisation of a theorem of Hatami, Janson and Szegedy on the entropy of graph limits. Finally, we raise a number of questions on geometric and analytic approaches to containers.


## KEYWORDS

containers, entropy, graph limits, hereditary bodies, rectilinear approximation

## 1 | INTRODUCTION

## 1.1 | Aims of the paper

In a major breakthrough 6 years ago now, Balogh-Morris-Samotij [5] and Saxton-Thomason [33] developed powerful theories of hypergraph containers. Given a hypergraph $H$ satisfying some smoothness assumptions they showed that there exists a small collection of almost independent sets whose subsets contain all the independent sets of $H$. A wide variety of problems in combinatorics are equivalent to estimating the number of independent sets in

[^0]various hypergraphs; the groundbreaking work of $[5,33]$ has thus seen an equally wide variety of applications, see, for example, the surveys [6, 7].

In this paper, our aim is to explore the implications of container theory beyond the discrete setting (which has hitherto been the main focus in applications) to the continuous setting, and to ask whether it is possible to obtain some form of containers going in the other direction, that is, starting from results in the continuous setting. We do this in two ways.

First of all, we relate hypergraph containers to rectilinear approximation of continuous bodies. Informally, we show that container theory implies the following: consider a sequence of bodies $\left(b_{n}\right)_{n \in \mathbb{N}}$, where $b_{n} \subseteq[0,1]^{d_{n}}$. Suppose this sequence is closed under certain projections (satisfying some simple, natural conditions). Then the bodies in the sequence can be finely approximated by a small number of boxes. This (informally stated) result, Theorem 1.12, allows us to apply container theory to functions from discrete structures to [0, 1]-for instance, we estimate (Theorem 3.2) the probability that a random function from the Boolean hypercube to $[0,1]$ is $c$-Lipschitz. We use Theorem 1.12 to obtain general volume estimates for hereditary bodies. Similarly to applications of container theory to counting problems, this requires certain supersaturation results, which in much of the literature are obtained in an ad hoc manner. One of our contributions in this paper is to obtain a general, widely applicable form of supersaturation under a natural assumption (which is satisfied in most examples that have been studied), leading to a very clean general statement, Theorem 1.16, for volume estimates. A key question arising from this part of the paper is whether our rectilinear approximation results could be obtained directly from purely geometric considerations: given a sequence of bodies, is being closed under some family of projections enough to ensure the existence of good rectilinear approximations without resorting to container machinery? Also could some (weak) form of container theorem be obtained from geometric approximation arguments?

Second, in what was the initial motivation of this work, we investigate links between hypergraph containers and the theory of graph limits. Via container theory, we prove a multicolour generalisation of a theorem of Hatami, Janson and Szegedy [22] on the entropy of graph limits (Theorem 4.2). Our work in this part of the paper leads us to two questions. Can one extend the Hatami-Janson-Szegedy theorem further to [0, 1]-decorated graph limits? This connects to a broader project of Lovász and Szegedy [28] on extending the theory of graph limits to limits of compact decorated graphs. Further, as above, is it possible to go in the other direction, and to derive some (weak) form of container theorem for graph properties from compactness results for graphons?

We note our work in this paper focuses exclusively on "thick" hereditary bodies, whose volume varies exponentially with the dimension, rather than "thin" bodies whose volume decays superexponentially. It is thus natural to ask whether one can obtain a set of streamlined general results similar to the ones we derive in this paper but for "thin" bodies. Also, our work suggests families of new Turán-type entropy maximisation problems. These, along with the questions raised above, are discussed in greater detail in Section 5.

## 1.2 | Background

The problem of estimating the number of members of a hereditary class of discrete objects and characterising their typical structure has a long and distinguished history, beginning with the work of Erdős, Kleitman and Rothschild [14] in the 1970s. The Alekseev-Bollobás-Thomason theorem [1, 2, 9] determined the asymptotics of the logarithm of the number of graphs on $n$ vertices in a hereditary property of graphs, while Alon, Balogh, Bollobás and Morris [4]
characterised their typical structure. Further Conlon and Gowers [12] and Schacht [35] obtained general transference results, which in particular implied sparse random analogues of extremal theorems for monotone properties of graphs (see the International Congress of Mathematicians [ICM] survey of Conlon [11] devoted to this topic).

In a major development in 2015, Balogh, Morris and Samotij [5] and independently Saxton and Thomason [33] developed powerful theories of hypergraph containers. Informally, they showed that-given some smoothness conditions-one may find in an $r$-uniform hypergraph $H$ on $n$ vertices a small ( $\operatorname{size} 2^{o\left(n^{2}\right)}$ ) collection $\mathcal{C}$ of almost independent sets (containing at most $o\left(n^{r}\right)$ edges), with the container property that every independent set $I$ in $H$ is contained inside some $C \in \mathcal{C}$. Provided one has a good understanding of the size and structure of the largest independent sets in $H$ (which is an extremal problem), one can use containers to estimate the number of independent sets in $H$ and characterise their typical structure, and to transfer such results to sparse random subhypergraphs of $H$. Since many well-studied hereditary properties of discrete structures correspond to the collection of independent sets in some suitably defined hypergraphs, the groundbreaking work of Balogh-Morris-Samotij and Saxton-Thomason has had an enormous number of applications, providing new, simplified proofs of many previous results as well as the resolution of many old conjectures. In the 6 years elapsed since their publication, the papers [5, 33] had amassed over 250 citations each. Among these let us note the works of greatest relevance to the present paper, namely, the work of Balogh and Wagner [7] showcasing the versatility of the container method, the papers of Terry [37] and Falgas-Ravry-O'Connell-Uzzell [18] on applications of containers to multicoloured discrete structures, and the ICM survey of Balogh, Morris and Samotij [6] devoted to hypergraph containers.

Following on [18], in which containers were adapted to the multicolour setting via random colouring models whose (discrete) entropy was used to count the underlying multicoloured structures, we shall in this paper use (continuous) entropy in combination with containers to estimate the volume of hereditary bodies. Entropy was introduced by Shannon [36] in a foundational paper on information theory; the use of entropy for counting (in the discrete setting) or making volume estimates (in the continuous setting) is a well-established technique in combinatorics, see, for example, the lectures of Galvin [20] on this topic. Mention should be made here of the recent and impressive results of Kozma, Meyerovitch, Peled and Samotij [24] who obtained a very fine approximation of the metric polytope (see the discussion in Section 3.2) via much more sophisticated and involved entropy techniques than the ones used in this paper.

One motivation for writing this paper was to better understand the potential links between container theory and the theory of graph limits. Giving an exposition of the latter theory is beyond the scope of this paper, and we refer the interested reader to the monograph of Lovász on the topic [26]. It suffices to say here that in the theory of (dense) graph limits one passes from the discrete world of graphs to the continuous world of graphons, which are symmetric measurable functions $W:[0,1]^{2} \rightarrow[0,1]$. One can then seek to recover many finitary results of graph theory in the limit world of graphons via analytic techniques (or in fact prove new results which can be exported back to the world of finite graphs). In a 2018 paper, Hatami, Janson and Szegedy [22] defined an entropy function for graphons, and used this entropy function to reformulate and give an alternative proof of the Alekseev-Bollobás-Thomason theorem in the graph limit setting. Lovász and Szegedy [28] began extending the theory of graph limits from ordinary graphs to graphs whose edges are decorated or coloured with elements from a compact set; further work in this direction was done more recently by Kunszenti-Kovács, Lovász and Szegedy [25], though (as we mention in Section 5) some parts of the theory are yet to be extended, such as the extraction of convergent subsequences from an arbitrary sequence of decorated graphs.

## 1.3 | Definitions and statement of our main results

Before we can state our main results, we need to introduce some basic notation, to recall definitions of set sequences equipped with embeddings (ssee-s) and to define a number of concepts related to $[0,1]$-decorations, entropy and hereditary properties of decorated ssee-s. Throughout the paper we shall use $|A|$ to denote the Lebesgue measure of $A$ when $A$ is a measurable subset of $[0,1]$. We also let $[n]$ denote the discrete interval $\{1,2, \ldots, n\}$. Finally, we shall use standard Landau notation: given functions $f, g: \mathbb{N} \rightarrow \mathbb{R}^{+}$, we write $f=o(g)$ for $\lim _{n \rightarrow \infty} f(n) / g(n)=0$ and $f=O(g)$ if there exists a constant $C>0$ such that $\limsup _{n \rightarrow \infty} f(n) / g(n) \leq C$. Further we write $f=\omega(g)$ for $g=o(f), f=\Omega(g)$ for $g=O(f)$.

### 1.3.1 | Set sequences equipped with embeddings

We begin by recalling the definition of a set sequence equipped with embeddings (ssee), and that of a good ssee, from the prequel [18] to this paper.

Definition 1.1 (ssee). A set sequence equipped with embeddings, or ssee, is a sequence $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$ of sets $V_{n}$, together with for every $N \leq n$ a collection $\binom{V_{n}}{V_{N}}$ of injections $\phi: V_{N} \rightarrow V_{n}$. We refer to the elements of $\binom{V_{n}}{V_{N}}$ as embeddings of $V_{N}$ into $V_{n}$.

Definition 1.2 (Intersecting embeddings). Let $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$ be an ssee. Let $N_{1}, N_{2} \leq n$. An i-intersecting embedding of ( $V_{N_{1}}, V_{N_{2}}$ ) into $V_{n}$ is a function $\phi: V_{N_{1}} \sqcup V_{N_{2}} \rightarrow V_{n}$ such that:
(i) the restriction of $\phi$ to $V_{N_{1}}$ lies in $\binom{V_{n}}{V_{N_{1}}}$, and the restriction of $\phi$ to $V_{N_{2}}$ lies in $\binom{V_{n}}{V_{N_{2}}}$;
(ii) $\left|\phi\left(V_{N_{1}}\right) \cap \phi\left(V_{N_{2}}\right)\right|=i$.

We denote by $I_{i}\left(\left(V_{N_{1}}, V_{N_{2}}\right), V_{n}\right)$ the number of $i$-intersecting embeddings of $\left(V_{N_{1}}, V_{N_{2}}\right)$ into $V_{n}$, and set

$$
I(N, n):=\sum_{1<i<\left|V_{N}\right|} I_{i}\left(\left(V_{N}, V_{N}\right), V_{n}\right) .
$$

Definition 1.3 (Good ssee). An ssee $\mathbf{V}$ is good if it satisfies the following conditions:
(i) $\left|V_{n}\right| \rightarrow \infty$ ("the sets in the sequence become large");
(ii) for all $N \in \mathbb{N}$ with $\left|V_{N}\right|>1,\left|\binom{V_{n}}{V_{N}}\right| \gg\left|V_{n}\right|$ ("on average, vertices in $V_{n}$ are contained in many embedded copies of $V_{N} "$ );
(iii) for all $N \in \mathbb{N}$ with $\left|V_{N}\right|>1,\left(\left|V_{n}\right| I(N, n)\right) /\left|\binom{V_{n}}{V_{N}}\right|^{2} \rightarrow 0$ as $n \rightarrow \infty$ ("most pairs of embeddings of $V_{N}$ into $V_{n}$ share at most one vertex").

The notion of an ssee covers a wide variety of well-studied structures: all of the following are examples of good ssee-s:

- $V_{n}$ is the edge-set of the complete graph on $n$ vertices $E\left(K_{n}\right)$, and $\binom{V_{n}}{V_{N}}$ is the collection of maps $E\left(K_{N}\right) \rightarrow E\left(K_{n}\right)$ corresponding to the collection of graph isomorphisms from $K_{N}$ into $K_{n}$;
- $V_{n}$ is $\{0,1\}^{n}$, the vertex-set of the $n$-dimensional hypercube $Q_{n}$, and $\binom{V_{n}}{V_{N}}$ is the collection of graph isomorphisms from $Q_{N}$ into $Q_{n}$;
- $V_{n}$ is [ $n$ ], the interval of the first $n$ natural numbers, and $\binom{V_{n}}{V_{N}}$ is the collection of all injections $\phi$ sending [ $N$ ] into an arithmetic progression of [ $n$ ] of length $N$, that is, $\phi: x \mapsto a+x d$, where $a, d$ are fixed nonnegative integers and $d>0$;
- $V_{n}$ is $\left(\mathbb{F}_{p}\right)^{n}$, where $p$ is a prime and $\mathbb{F}_{p}$ is the finite field with $p$ elements, and $\binom{V_{n}}{V_{N}}$ is the collection of shifts of injective additive homomorphisms from $\left(\mathbb{F}_{p}\right)^{N}$ to $\left(\mathbb{F}_{p}\right)^{N}$-subgroups of $\left(\mathbb{F}_{p}\right)^{n}$ (i.e., the collection of all maps $\phi: \mathbf{a}+\psi(\mathbf{x})$, where $\mathbf{a} \in\left(\mathbb{F}_{p}\right)^{n}, \psi:\left(\mathbb{F}_{p}\right)^{N} \rightarrow\left(\mathbb{F}_{p}\right)^{n}$ is injective and satisfies $\psi(\mathbf{x}+\mathbf{y})=\psi(\mathbf{x})+\psi(\mathbf{y})$; the image of $\left(\mathbb{F}_{p}\right)^{N}$ under $\phi$ is thus a coset of an $\left(\mathbb{F}_{p}\right)^{N}$-subgroup of $\left.\left(\mathbb{F}_{p}\right)^{n}\right)$;
- somewhat similar to the example above, $V_{n}$ is $\left(\mathbb{F}_{2}\right)^{n} \backslash\{\mathbf{0}\}$, and $\binom{V_{n}}{V_{N}}$ is the set of linear isomorphisms from $\left(\mathbb{F}_{2}\right)^{N}$ to $n$-dimensional linear subspaces of $\left(\mathbb{F}_{2}\right)^{n}$, restricted to $\left(\mathbb{F}_{2}\right)^{N} \backslash\{\mathbf{0}\}$ (this example corresponds to simple binary matroids $M:\left(\mathbb{F}_{2}\right)^{n} \backslash\{\mathbf{0}\} \rightarrow\{0,1\}$, whose hereditary properties were recently investigated by Grosser, Hatami, Nelson and Norin [21]);
- $V_{n}$ is the collection $\mathcal{P}([n])$ of all subsets of $[n]$, viewed as a poset under the subset relation, and $\binom{V_{n}}{V_{N}}$ is the collection of all injective poset homomorphisms from $\mathcal{P}([N])$ into $\mathcal{P}([n])$ (these homomorphisms can be counted using [16, Theorem 4.1], from which properties (ii) and (iii) follow easily).


### 1.3.2 | [0, 1]-Decorations, entropy, hereditary properties

We now generalise a number of definitions from [18] to the setting of [0, 1]-decorated ssee-s. In what follows, we write $\left(v_{i}\right)_{i \in I}$ to denote a vector $v$ whose coordinates are labelled with elements of some index set $I$, and refer to such vectors as $I$-indexed vectors. Given a set $S$, we also write $S^{I}$ for the collection of $I$-indexed vectors all of whose coordinates take values in $S$, that is, for the $I$-indexed Cartesian product $\prod_{i \in I} S$.

Definition $1.4[0,1](-$ Decorated sets and set sequences). Given a set $V$, a $[0,1]$ decoration of $V$ is an element $c \in[0,1]^{V}$, that is, a function $c: V \rightarrow[0,1]$. Given a set sequence $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$, a $[0,1]$-decoration of $\mathbf{V}$ is a sequence $\mathbf{c}=\left(c_{n}\right)_{n \in \mathbb{N}}$, where for each $n \in \mathbb{N}, c_{n}$ is a [0, 1]-decoration of $V_{n}$. If $\mathbf{V}$ is an ssee, we call $\mathbf{c}$ a $[0,1]$-decorated ssee.
$\mathrm{A}[0,1]$-decorated set is just a function from the set into $[0,1]$. Our interest in this paper is the extent to which such hereditary families of such functions can be approximated by (rectilinear) boxes, which we define below.

Definition 1.5 (Boxes and cylinders). A box in $V_{n}$ is a Cartesian product of the form

$$
b:=\prod_{i \in V_{n}} A_{i},
$$

where for each $i, A_{i}$ is a measurable subset of $[0,1]$. (Thus $b$ is a collection of $V_{n}$-indexed vectors.) We write $\operatorname{Box}\left(V_{n}\right)$ for the collection of all boxes in $V_{n}$. A $d$-cylinder is a box where all but at most $d$ of the $A_{i}$ are equal to $[0,1]$.

A cylinder or box is said to be simple if for every $i \in V_{n}$, its projection onto coordinate $i$ is a finite union of intervals. Further, a simple box is called $k$-rational if each of these intervals is of the form $\left[\frac{x}{k}, \frac{y}{k}\right]$, for some integers $0 \leq x<y \leq k$.
Definition 1.6 (Volume, entropy and density). Given a measurable body $b \subseteq[0,1]^{V_{n}}$, we denote its Lebesgue measure by vol $(b)$. So, for instance, if $b=\prod_{i \in V_{n}} A_{i}$ is a box, its volume is $\operatorname{vol}(b)=\prod_{i \in V_{n}}\left|A_{i}\right|$.

We also consider the volume of lower-dimensional bodies obtained from $b$ by fixing the coordinates inside some subset $I \subseteq V_{n}$; we then denote the corresponding $\left|V_{n} \backslash I\right|$-dimensional volume (with respect to the Lebesgue measure) by vol ${l_{n \backslash} \backslash .}$. So, for instance, for $v_{0} \in V_{n}$

$$
\operatorname{vol}_{V_{n}\left\{\left\{v_{0}\right\}\right.}\left(\left\{c \in b: c_{v_{0}}=1 / 2\right\}\right)
$$

denotes the volume of the ( $\left|V_{n}\right|-1$ )-dimensional object obtained by taking the intersection of $b$ with the hyperplane $\left\{x \in \mathbb{R}^{V_{n}}: x_{v_{0}}=1 / 2\right\}$.

Further, for a measurable body $b \subseteq[0,1]^{V_{n}}$, we define the entropy of $b$ as

$$
\operatorname{Ent}(b):=-\log \operatorname{vol}(b) .
$$

Finally, the density of $b$ is

$$
d(b):=\operatorname{vol}(b)^{1 / / V_{n}} .
$$

Observe that this latter quantity is an element of $[0,1]$, and that $d(b)=e^{-\operatorname{Ent}(b) /\left|V_{n}\right|}$.
We now turn to the problem of defining what we mean by a hereditary family of functions on an ssee. To do this, we first define a notion of projection inherited from the embeddings associated with the ssee.

Definition 1.7 (Projections, lifts and shadows). Let $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$ be an ssee. Given an embedding $\phi \in\binom{V_{n}}{V_{N}}$ and a measurable body $b \subseteq[0,1]^{V_{n}}$, we denote by $b_{\downarrow \phi}$ the measurable subset of $[0,1]^{V_{V}}$ given by

$$
b_{\downarrow \phi}:=\left\{c \in[0,1]^{V_{N}}: \operatorname{vol}_{V_{n} \backslash \phi\left(V_{N}\right)}\left(\left\{\tilde{c} \in b: \tilde{c}_{\phi(i)}=c_{i} \forall i \in V_{N}\right\}\right)>0\right\} .
$$

So, for example, given a box $b=\prod_{i \in V_{n}} A_{i}$, we have

$$
b_{\downarrow}:= \begin{cases}\prod_{i \in V_{N}} A_{\phi(i)} & \text { if } \operatorname{vol}\left(A_{j}\right)>0 \quad \forall j \in V_{n} \backslash \phi\left(V_{N}\right) \\ \prod_{i \in V_{N}} \varnothing & \text { otherwise } .\end{cases}
$$

We call $b_{\downarrow \phi}$ the (strict) $\phi$-projection of $b$.
Conversely, given an embedding $\phi \in\binom{V_{n}}{V_{N}}$ and a measurable body $b \subseteq[0,1]^{V_{N}}$, we denote by $b_{\uparrow \phi}$ the measurable subset of $[0,1]^{V_{n}}$ induced by $\phi$, namely,

$$
b_{\uparrow \phi}:=\left\{c \in[0,1]^{V_{n}}: \exists \tilde{c} \in b \text { such that } \forall i \in V_{N}, \tilde{c}_{i}=c_{\phi(i)}\right\} .
$$

So, for example, given a box $b=\prod_{i \in V_{N}} A_{i}$, we have $b_{\uparrow \phi}:=\prod_{i \in V_{n}} B_{i}$, where

$$
B_{i}:= \begin{cases}A_{\phi^{-1}(i)} & \text { if } i \in \phi\left(V_{N}\right), \\ {[0,1]} & \text { otherwise }\end{cases}
$$

We call $b_{\uparrow \phi}$ the $\phi$-lift of $b$.
Given $b \in[0,1]^{V_{n}}$ and $N \leq n$, we define the lower shadow of $b$ in $[0,1]^{V_{N}}$ by

$$
\partial_{V_{N}}^{-}(b):=\bigcup_{\phi \in\binom{V_{n}}{V_{N}}} b_{\emptyset \phi} .
$$

Similarly, given $b \in[0,1]^{V_{N}}$ and $n \geq N$, we define the upper shadow of $b$ in $[0,1]^{V_{n}}$ by

$$
\partial_{V_{n}}^{+}(b):=\underset{\phi \in\binom{V_{n}}{V_{N}}}{\bigcup} b_{\uparrow \phi} .
$$

Observe that if $b$ is a box in $V_{n}$ and $\phi \in\binom{V_{n}}{V_{N}}$, then $b_{\downarrow \phi} \in \operatorname{Box}\left(V_{N}\right)$. Conversely, if $b$ is a box in $V_{N}$ and $\phi \in\binom{V_{n}}{V_{N}}$, then $b_{\uparrow \phi}$ is a box (in fact a $\left|V_{N}\right|$-cylinder) in $V_{n}$. Also, for any body $b \subseteq[0,1]^{V_{N}}$ and any embeddings $\phi \in\binom{V_{n}}{V_{N}}, \psi \in\binom{V_{N}}{V_{n^{\prime}}}$ we have the relations

$$
\left(b_{\uparrow \phi}\right)_{\downarrow \phi}=b \quad \text { and } \quad \operatorname{vol}\left(b \backslash\left(b_{\downarrow \psi}\right)_{\uparrow \psi}\right)=0 .
$$

Definition 1.8 (Properties and hereditary properties). Let $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$ be an ssee. A $[0,1]$-decoration property of $\mathbf{V}$ is a sequence $\mathcal{P}=\left(\mathcal{P}_{n}\right)_{n \in \mathbb{N}}$, where $\mathcal{P}_{n}$ is a measurable subset of $[0,1]^{V_{n}}$. A $[0,1]$-decoration property is hereditary if for all $n \geq N$, the upper shadow of $\left([0,1]^{V_{N}} \backslash \mathcal{P}_{N}\right)$ in $[0,1]^{V_{n}}$ is a subset of $\left([0,1]^{V_{n}} \backslash \mathcal{P}_{n}\right)$.

In other words, a [0, 1]-decoration property $\mathcal{P}$ of $\mathbf{V}$ is hereditary if its complement is closed under taking upper shadows $/ \boldsymbol{\phi}$-lifts. In particular, this implies that $\mathcal{P}$ itself is closed under taking lower shadows/ $\phi$-projections. (Note however that the converse fails: a property being closed under taking lower shadows does not imply its complement is closed under taking upper shadows.)

As an illustrative example, observe that our abstract definition above generalises the graphtheoretic notion of a hereditary property. Indeed let $\left(\mathcal{G}_{n}\right)_{n \in \mathbb{N}}$ be a sequence of subgraphs of $K_{n}$, the complete graph on $n$ vertices, that is closed under taking induced subgraphs (e.g., one could take $\mathcal{G}_{n}$ to be the collection of all triangle-free subgraphs of $K_{n}$, or of all the subgraphs of $K_{n}$ containing no induced cycle of length 5 , say). Letting $V_{n}$ denote the edge-set of $K_{n}$, we can
encode a graph $G \in \mathcal{G}_{n}$ as a box $b_{G}=\prod_{e \in E\left(K_{n}\right)} A_{e}$ by letting $A_{e}=\left[\frac{1}{2}, 1\right]$ if $e \in G$ and setting $A_{e}=\left[0, \frac{1}{2}\right)$ otherwise. Letting $\mathcal{P}_{n}:=\bigcup_{G \in \mathcal{G}_{n}} b_{G}$, we have that the [0, 1]-decoration property $\mathcal{P}_{n}$ of the ssee $\mathcal{V}=\left(V_{n}\right)_{n \in \mathcal{N}}$ is hereditary in the sense of Definition 1.8.

One class of hereditary [0, 1]-decoration properties will be of particular interest to us in this paper.

Definition 1.9 (Forbidden projections). Let $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$ be an ssee. Let $b \subseteq[0,1]^{V_{N}}$. For $n \geq N$, we say $b^{\prime} \subseteq[0,1]^{V_{n}}$ is $b$-free if $\partial_{V_{n}}^{+}(b) \cap b^{\prime}=\varnothing$. We denote by Forb $(b)$ the hereditary [0, 1]-decoration property of being $b$-free, that is, for all $n \geq N$,

$$
\operatorname{Forb}(b)_{n}=[0,1]^{V_{n}} \backslash \partial_{V_{n}}^{+}(b)
$$

Definition 1.10 (Extremal entropy). Let $\mathbf{V}=\left(V_{n}\right)_{n \in \mathbb{N}}$ be an ssee, and let $\mathcal{P}=\left(\mathcal{P}_{n}\right)_{n \in \mathbb{N}}$ be a $[0,1]$-decoration property of $\mathbf{V}$. The extremal entropy of $\mathcal{P}$ relative to $\mathbf{V}$ is

$$
\operatorname{ex}(\mathbf{V}, \mathcal{P})_{n}=\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right):=\inf \left\{\operatorname{Ent}(b): b \in \operatorname{Box}\left(V_{n}\right), \operatorname{vol}\left(b \backslash \mathcal{P}_{n}\right)=0\right\} .
$$

Thus $\exp \left(-\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right)\right)$ is precisely the volume of the largest box in $V_{n}$ which (up to a zeromeasure set) is contained inside $P_{n}$.

### 1.3.3 | Rectilinear approximation and volume estimates for hereditary bodies

Let $\mathbf{V}$ be a good ssee and let $n, N \in \mathbb{N}$. Given a function $f:\binom{V_{n}}{V_{N}} \rightarrow \mathbb{R}$, we write $\mathbb{E}_{\phi} f(\phi)$ for the expected value of $f(\phi)$ over $\phi \in\binom{V_{n}}{V_{N}}$ chosen uniformly at random.

Our first result is a geometric approximation property for a hereditary body.
Theorem 1.11. Let $\mathbf{V}$ be a good ssee. Let $\mathcal{F}$ be a nonempty finite family of simple boxes in $[0,1]^{V_{N}}$, for some $N \in \mathbb{N}$. Set $b=\bigcup_{b_{f} \in \mathcal{F}} b_{f}$ and $\mathcal{P}=\operatorname{Forb}(b)$. Then for every $\varepsilon>0$, there exists $n_{0}>0$ such that for any $n \geq n_{0}$ there exists a collection $\mathcal{C}$ of simple boxes in $[0,1]^{V_{n}}$ satisfying:
(i) $\mathcal{P}_{n} \subseteq \bigcup_{c \in \mathcal{C}} c$;
(ii) for every $c \in \mathcal{C}, \mathbb{E}_{\phi} \operatorname{vol}\left(c_{\downarrow \phi} \cap b\right)<\varepsilon$;
(iii) $\mid \mathcal{C l} \leq e^{\varepsilon\left|V_{n}\right|}$.

In other words there exists a small (property [iii]) collection of simple boxes such that their union contains the body $\mathcal{P}_{n}$ (property [i]). Further, each of them has a lower shadow almost disjoint from $b$ (property [ii])—so these boxes "almost" lie in $\mathcal{P}_{n}$. The union of these boxes is thus a "good" approximation for $\mathcal{P}_{n}$.

It is worth noting that of course any measurable body can be finely approximated by a collection of simple boxes-the power of the container theory of Balogh-Morris-Samotij and Saxton-Thomason is the bound (iii) they give on the number of boxes required.

Building on Theorem 1.11, we prove:
Theorem 1.12. Let $\mathbf{V}$ be a good ssee. Let $\mathcal{P}$ be a hereditary property of [0, 1]-decorations of $\mathbf{V}$ and let $N \in \mathbb{N}$. Then for every $\varepsilon>0$, there exists an integer $n_{0}>N$ such that for any $n \geq n_{0}$ there exists a collection $\mathcal{C}$ of simple boxes in $[0,1]^{V_{n}}$ satisfying:
(i) $\operatorname{vol}\left(\mathcal{P}_{n} \backslash \bigcup_{c \in \mathcal{C}} c\right)=0$;
(ii) for every $c \in \mathcal{C}, \mathbb{E}_{\phi} \operatorname{vol}\left(\mathcal{c}_{\downarrow \phi} \backslash \mathcal{P}_{N}\right)<\varepsilon$;
(iii) $\mid \mathcal{C l} \leq e^{\varepsilon\left|V_{n}\right|}$.

We remark that we cannot replace (i) by a containment condition $\mathcal{P}_{n} \subseteq \bigcup_{c \in \mathcal{C}} \mathcal{C}$, and that we must allow for an exceptional zero-measure set not covered by the simple boxes in $\mathcal{C}$. For instance, suppose $\mathcal{P}$ consisted of all decorations $x \in[0,1]^{V_{n}}$ with $x_{i} \in\left[0, \frac{1}{2}\right] \cup \mathbb{Q}$. Then clearly we cannot both cover all of $\mathcal{P}_{n}$ with simple boxes and still achieve (ii).

Observe also that condition (ii) may be interpreted as follows: suppose we take a point $\mathbf{x}$ chosen uniformly at random from $c$ and an embedding $\phi$ uniformly at random from $\binom{V_{n}}{V_{N}}$. This defines a random point $\mathbf{y} \in[0,1]^{V_{N}}$, by setting $y_{v}=x_{\phi(v)}$ for all $v \in V_{N}$. Then condition (ii) is saying that the probability $\mathbf{y}$ fails to be in $\mathcal{P}_{N}$ is small (at most $\varepsilon$ ).

Provided we have (a) a limiting density and (b) supersaturation for a [0, 1]-decoration property in a good ssee, the container/geometric approximation theorem, Theorem 1.12, immediately implies a volume estimate for the $V_{n}$-dimensional body $\mathcal{P}_{n}$, namely:

Corollary 1.13. Let $\mathbf{V}$ be a good ssee. Let $\mathcal{P}$ be a hereditary property of [0, 1]-decorations of V. Suppose in addition that the following hold:
(a) $\pi(\mathcal{P}):=\lim _{n \rightarrow \infty} \operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right) /\left|V_{n}\right|$ exists;
(b) for all $\varepsilon>0$, there exist $\eta>0$ and positive integers $N \leq n_{0}$ such that if $n \geq n_{0}$ then for every $b \in \operatorname{Box}\left(V_{n}\right)$, if $\mathbb{E}_{\phi} \operatorname{vol}\left(b_{\downarrow \phi} \backslash \mathcal{P}_{N}\right)<\eta$ then $\operatorname{Ent}(b)>(\pi(\mathcal{P})-\varepsilon)\left|V_{n}\right|$.

Then

$$
\operatorname{vol}\left(\mathcal{P}_{n}\right)=e^{-(\pi(\mathcal{P})+o(1))\left|V_{n}\right|}
$$

A natural question is whether we can give a simple criterion for satisfying assumptions (a) and (b). In past applications, this has mostly been dealt with in an ad hoc manner. For example, in [18], it was proved that these two assumptions were satisfied for vertex $k$ colourings and edge $k$-colourings of both complete hypergraphs and hypercube graphs, but each case required a separate proof. One of our contributions in this paper is a very simple criterion on the family of embeddings which is sufficient to ensure (a) and (b) are satisfied. This criterion immediately applies to a wide class of structures, yielding volume estimate results of great generality.

Definition 1.14. An ssee $\mathbf{V}$ is homogeneous if for every $n \geq N$, every $x \in V_{n}$ is contained in the same, strictly positive number of embeddings $\phi\left(V_{N}\right), \phi \in\binom{V_{n}}{v_{N}}$.

Proposition 1.15. Suppose $\mathbf{V}$ is a good homogeneous ssee. Then for every hereditary property $\mathcal{P}$ of $[0,1]$-decorations of $\mathbf{V}$, the sequence $\frac{\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right)}{\left|V_{n}\right|}$ is nondecreasing in $\mathbb{R}_{\geq 0}$. In particular, this sequence either converges to a limit $\pi(\mathcal{P})$ or tends to infinity as $n \rightarrow \infty$.

Theorem 1.16. Suppose $\mathbf{V}$ is a good homogeneous ssee. Then for every hereditary property $\mathcal{P}$ of $[0,1]$-decorations of $\mathbf{V}$, either

$$
\lim _{n \rightarrow \infty} \frac{\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right)}{\left|V_{n}\right|}=\pi(\mathcal{P})
$$

exists and

$$
\operatorname{vol}\left(\mathcal{P}_{n}\right)=e^{-(\pi(\mathcal{P})+o(1))\left|V_{n}\right|}
$$

or $\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right) /\left|V_{n}\right| \rightarrow \infty$ as $n \rightarrow \infty$ and the volume of $\mathcal{P}_{n}$ decays superexponentially in $\left|V_{n}\right|, \operatorname{vol}\left(\mathcal{P}_{n}\right) \leq e^{-\omega\left(\left|V_{n}\right|\right)}$.

## 1.4 | Structure of the paper

We prove our main results in Section 2. In Section 3 we give simple applications of our results to Lipschitz functions on the hypercube, the metric polytope, and weighted graphs, while in Section 4 we use our results to derive a multicolour/decorated version of an "entropy of graph limits" theorem of Hatami, Janson and Szegedy [22] (which was one of our initial motivations for undertaking this project). We end the paper in Section 5 with a discussion of various questions arising from our work.

## 2 | PROOFS OF THE THEORETICAL RESULTS

For $k \in \mathbb{N}$, let $[k]$ denote the set $\{1,2, \ldots, k\}$. One of the main results of the prequel to this paper [18] was a container theorem for [ $k$ ]-decorated ssee-s (which itself was obtained as a consequence of a simple container theorem of Saxton and Thomason [34]). To state it we must recall a few definitions.

A $k$-colouring template for $V_{n}$ is a function $t$ assigning to each element of $V_{n}$ a nonempty subset of $[k]$. A realisation of $t$ is a colouring $c: V_{n} \rightarrow[k]$ with the property that $c(x) \in t(x)$ for all $x \in V_{n}$. The set of all realisations of $t$ is denoted by $\langle t\rangle$. Given a collection $\mathcal{F}$ of $k$-colourings of $V_{N}$, denote by $\operatorname{Forb}_{\mathbf{V}}(\mathcal{F})$ the hereditary property of $k$-colourings of $\mathbf{V}$ not containing an embedding of a colouring in $\mathcal{F}$, that is,

$$
\begin{equation*}
\operatorname{Forb}_{\mathbf{v}}(\mathcal{F})_{n}:=\left\{c \in[k]^{V_{n}}: \forall \phi \in\binom{V_{n}}{V_{N}}, c_{\downarrow \phi} \notin \mathcal{F}\right\} \tag{1}
\end{equation*}
$$

where in the $[k]$-decorated context $c_{\downarrow \phi}$ denotes the $k$-colouring of $V_{N}$ induced by $\phi, c_{\downarrow \phi}: v \mapsto c(\phi(v))$.

Theorem 2.1 (Falgas-Ravry et al. [18, Theorem 3.18]). Let $\mathbf{V}$ be a good ssee, and let $k, N \in \mathbb{N}$. Let $\mathcal{F}$ be a nonempty subset of $[k]^{V_{N}}$ and $\mathcal{P}=\operatorname{Forb}_{\mathbf{v}}(\mathcal{F})$. For any $\varepsilon>0$, there exists $n_{0}>0$ such that for any $n \geq n_{0}$ there exists a collection $\mathcal{T}_{n}$ of $k$-colouring templates for $V_{n}$ satisfying:
(a) $\mathcal{P}_{n} \subseteq \bigcup_{t \in \mathcal{T}_{n}}\langle t\rangle$;
(b) for each template $t \in \mathcal{T}_{n}$, there are at most $\varepsilon\left|\binom{V_{n}}{V_{N}}\right|$ pairs $(\phi, c)$ with $\phi \in\binom{V_{n}}{V_{N}}, c \in \mathcal{F}$ and $c \in\left\langle t_{\downarrow \phi}\right\rangle$;
(c) $\left|\mathcal{T}_{n}\right| \leq k^{\varepsilon\left|V_{n}\right|}$.

Our strategy to prove the theorems in Section 1.3 .3 is to use compactness: as $[0,1]$ is bounded, we can approximate measurable sets by finite unions of intervals of the form $\left[\frac{i-1}{k}, \frac{i}{k}\right]$, where $k \in \mathbb{N}$ is some large constant and $i \in[k]$. There we can apply the container theorem for colourings by discrete finite sets, Theorem 2.1. Provided we are careful with our approximations (and, crucially, that we are using the right definitions), we are able to transfer the container results from the discrete to the continuous setting.

Proof of Theorem 1.11. Fix $\varepsilon>0$. There exists $k \in \mathbb{N}$ such that there exists a finite union of $k$-rational simple boxes $\tilde{b}$ such that (1) $\tilde{b} \subseteq b$ and (2) $\operatorname{vol}(b \backslash \tilde{b})<\varepsilon / 2$. We can now pass to the discrete setting and consider $[k]$-colourings of $V_{N}$ as a proxy for $k$-rational simple boxes in $[0,1]^{V_{N}}$. Let

$$
\mathcal{F}:=\left\{c \in[k]^{V_{N}}: \prod_{v \in V_{N}}\left[\frac{c_{v}-1}{k}, \frac{c_{v}}{k}\right] \subseteq \tilde{b}\right\}
$$

be the family of $k$-colourings of $V_{N}$ corresponding to $\tilde{b}$. Apply Theorem 2.1 to $\mathcal{F}$ with parameters $k, N$ and $\varepsilon^{\prime}=\min (\varepsilon / \log k, \varepsilon / 2)$. Let $n_{0} \in \mathbb{N}$ be such that for all $n \geq n_{0}$ conclusions (a)-(c) from Theorem 2.1 hold. Let $\mathcal{T}_{n}$ be the family of templates whose existence is guaranteed by the theorem. For each $t \in \mathcal{T}_{n}$, we define a box $c^{t}$ given by

$$
c^{t}:=\prod_{v \in V_{n}}\left(\bigcup_{i \in t(v)}\left[\frac{i-1}{k}, \frac{i}{k}\right]\right)
$$

Let $\mathcal{C}$ denote the collection of boxes thus obtained. Property (a) from Theorem 2.1 implies

$$
\begin{equation*}
\mathcal{P}_{n}=\left([0,1]^{V_{n}} \backslash\left(\partial_{V_{n}}^{+}(b)\right)\right) \subseteq\left([0,1]^{V_{n}} \backslash\left(\partial_{V_{n}}^{+}(\tilde{b})\right)\right) \subseteq \bigcup_{t \in \mathcal{T}_{n}} c^{t}=\bigcup_{c \in \mathcal{C}} c . \tag{2}
\end{equation*}
$$

Note that the second containment relation in (2) follows since by definition, the (continuous) upper shadow $\partial_{V_{n}}^{+}(\tilde{b})$ when $\tilde{b}$ is a union of $k$-rational simple boxes corresponds precisely to the (discrete set of) $c \in[k]^{V_{n}}$ which are not in $\operatorname{Forb}_{\mathbf{v}}(\mathcal{F})_{n}$ (see Equation 1). This is what motivated our choice of $\mathcal{F}$ above.

Further, property (b) entails that if one fixes $t \in \mathcal{T}_{n}$ and picks $\phi \in\binom{V_{n}}{V_{M}}$ uniformly at random, there is at most an $\varepsilon^{\prime}$-chance that $\left(c^{t}\right)_{l \phi}$ intersects the box $\tilde{b}$ in a set with nonzero measure. Thus

$$
\begin{align*}
\mathbb{E}_{\phi} \operatorname{vol}\left(\left(c^{t}\right)_{\downarrow \phi} \cap b\right) & =\mathbb{E}_{\phi} \operatorname{vol}\left(\left(c^{t}\right)_{\downarrow \phi} \cap(b \backslash \tilde{b})\right)+\mathbb{E}_{\phi} \operatorname{vol}\left(\left(c^{t}\right)_{\downarrow \phi} \cap \tilde{b}\right) \\
& \leq \operatorname{vol}(b \backslash \tilde{b})+\varepsilon^{\prime}<\varepsilon . \tag{3}
\end{align*}
$$

Finally, property (c) gives

$$
\begin{equation*}
|\mathcal{C}|=\left|\mathcal{T}_{n}\right| \leq k^{\varepsilon^{\prime}\left|V_{n}\right|} \leq e^{\varepsilon\left|V_{n}\right|} . \tag{4}
\end{equation*}
$$

Together, (2)-(4) show properties (i)-(iii) in the statement of Theorem 1.11 are satisfied as claimed, concluding the proof.

Proof of Theorem 1.12. Fix $\varepsilon>0$. Let $b:=[0,1]^{V_{N}} \backslash \mathcal{P}_{N}$. Since $\mathcal{P}$ is a measurable set, there exists a finite union of simple boxes $\tilde{b}$ such that (1) $\operatorname{vol}(\tilde{b} \backslash b)=0$ (i.e., up to a zeromeasure set, $\tilde{b} \subseteq b$ ) and (2) $\operatorname{vol}(b \backslash \tilde{b})<\varepsilon / 2$, by the definition of the Lebesgue measure. Apply Theorem 1.11 to $\mathcal{Q}:=\operatorname{Forb}(\tilde{b})$ with parameter $\varepsilon / 2$, and let $\mathcal{C}$ be the resulting family of simple boxes.

As $\mathcal{P}$ is hereditary we have $\partial_{V_{n}}^{+}(b) \subseteq[0,1]^{V_{n}} \backslash \mathcal{P}_{n}$, which implies that

$$
\operatorname{vol}\left(\mathcal{P}_{n} \backslash \mathcal{Q}_{n}\right)=\operatorname{vol}\left(\mathcal{P}_{n} \cap \partial_{V_{n}}^{+}(\tilde{b})\right) \leq \operatorname{vol}\left(\mathcal{P}_{n} \cap \partial_{V_{n}}^{+}(b)\right)+\operatorname{vol}\left(\partial_{V_{n}}^{+}(\tilde{b}) \backslash \partial_{V_{n}}^{+}(b)\right)=0
$$

It follows that

$$
\operatorname{vol}\left(\mathcal{P}_{n} \backslash \bigcup_{c \in \mathcal{C}} c\right) \leq \operatorname{vol}\left(\mathcal{P}_{n} \backslash \mathcal{Q}_{n}\right)+\operatorname{vol}\left(\mathcal{Q}_{n} \backslash \bigcup_{c \in \mathcal{C}} c\right)=0
$$

Further, we have

$$
\operatorname{vol}\left(\mathcal{Q}_{N} \backslash \mathcal{P}_{N}\right)=\operatorname{vol}\left([0,1]^{V_{N}} \backslash\left(\partial_{V_{N}}^{+}(\tilde{b}) \cup \mathcal{P}_{N}\right)\right)=\operatorname{vol}(b \backslash \tilde{b})<\frac{\varepsilon}{2}
$$

and so for every $c \in \mathcal{C}$ we have

$$
\mathbb{E}_{\phi} \operatorname{vol}\left(c_{\downarrow \phi} \backslash \mathcal{P}_{N}\right) \leq \mathbb{E}_{\phi}\left(\operatorname{vol}\left(c_{\downarrow \phi} \backslash \mathcal{Q}_{N}\right)+\operatorname{vol}\left(\mathcal{Q}_{N} \backslash \mathcal{P}_{N}\right)\right)<\frac{\varepsilon}{2}+\frac{\varepsilon}{2}=\varepsilon
$$

Finally, $\left|\mathcal{C}_{n}\right| \leq e^{\varepsilon\left|V_{n}\right| / 2}$. Thus $\mathcal{C}$ satisfies the properties (i)-(iii) claimed by Theorem 1.12, as desired.

Proof of Corollary 1.13. Fix $\varepsilon>0$. Let $\eta>0$ and $N, n_{0} \in \mathbb{N}$ be the constants guaranteed by assumption (b). Applying Theorem 1.12 to $\mathcal{P}$ with parameter $\delta=\min (\varepsilon, \eta)$, we find $n_{1} \geq n_{0}$ such that for all $n \geq n_{1}$ there exists a collection $\mathcal{C}$ of simple boxes in $[0,1]^{V_{n}}$ such
that (i) up to a zero-volume set, $\mathcal{P}_{n}$ is contained inside $\bigcup_{c \in \mathcal{C}} c$, (ii) for every $c \in \mathcal{C}, \mathbb{E}_{\phi} \operatorname{vol}\left(c_{\mid \phi} \backslash \mathcal{P}_{N}\right)<\delta$ and (iii) $\mid \mathcal{C l} \leq e^{\delta\left|V_{n}\right|}$.

By assumption (b) and our choice of $\delta$ and $n_{1}$, this implies that for every $c \in \mathcal{C}$ we have $\operatorname{Ent}(c) \geq(\pi(\mathcal{P})-\varepsilon)\left|V_{n}\right|$. Now (i) and (iii) allow us to bound the volume of $\mathcal{P}_{n}$ :

$$
\begin{aligned}
\operatorname{vol}\left(\mathcal{P}_{n}\right) & \leq \operatorname{vol}\left(\bigcup_{c \in \mathcal{C}} c\right)=\sum_{c \in \mathcal{C}} e^{-\operatorname{Ent}(c)} \leq \mid \mathcal{C l} e^{-(\pi(\mathcal{P})-\varepsilon)\left|V_{n}\right|} \leq e^{\delta\left|V_{n}\right|} e^{-(\pi(\mathcal{P})-\varepsilon)\left|V_{n}\right|} \\
& \leq e^{-(\pi(\mathcal{P})-2 \varepsilon)\left|V_{n}\right|} .
\end{aligned}
$$

Since $\varepsilon>0$ was arbitrary, the theorem follows.
Proof of Proposition 1.15. Assume that $\mathbf{V}$ is a good, homogeneous ssee, and let $\mathcal{P}$ be a hereditary property of $[0,1]$-decorations of $\mathbf{V}$. Set $x_{n}=\frac{\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right)}{\left|V_{n}\right|}$. Let $b \in \operatorname{Box}\left(V_{n+1}\right)$ be a box with $\operatorname{vol}\left(b \backslash \mathcal{P}_{n+1}\right)=0$ and $\operatorname{Ent}(b)=\operatorname{ex}\left(V_{n+1}, \mathcal{P}_{n+1}\right)$. By homogeneity of the ssee $\mathbf{V}$, each coordinate in $V_{n+1}$ is counted in the same nonzero number $k=\left|\binom{V_{n+1}}{V_{n}}\right| \frac{\left|V_{n}\right|}{\left|V_{n+1}\right|}$ of projections $\phi\left(V_{n}\right)$, so that the family $\left\{\phi\left(V_{n}\right): \phi \in\binom{V_{n}+1}{V_{n}}\right\}$ constitutes a $k$-uniform cover of $V_{n+1}$. Since $b$ is a box, it follows that

$$
\begin{equation*}
\operatorname{vol}(b)\left|\binom{V_{n+1}}{V_{n}}\right| \frac{\left|V_{V^{\prime}}\right|}{\left|V_{n+1}\right|}=\prod_{\phi \in\binom{V_{n+1}}{V_{n}}} \operatorname{vol}\left(b_{\downarrow \phi}\right) . \tag{5}
\end{equation*}
$$

Since $\mathcal{P}$ is hereditary and $\operatorname{vol}\left(b \backslash \mathcal{P}_{n+1}\right)=0$, for every $\phi \in\binom{V_{n+1}}{V_{n}}$ the $\phi$-projection $b_{\downarrow \phi}$ is a box in $V_{n}$ satisfying $\operatorname{vol}\left(b_{\downarrow \phi} \backslash \mathcal{P}_{n}\right)=0$. In particular, we must have

$$
\operatorname{vol}\left(b_{\downarrow \phi}\right) \leq e^{-\operatorname{ex}\left(V_{n}, \mathcal{P}\right)}=e^{-x_{n}\left|V_{n}\right|}
$$

Combining this with (5), we have

$$
\left.e^{-x_{n+1} \mid}\left|\binom{V_{n+1}}{V_{n}}\right| V_{n}|\quad=\operatorname{vol}(b)|\binom{V_{n+1}}{V_{n}} \right\rvert\, \frac{\left|V_{n}\right|}{\left|V_{n+1}\right|} \leq e^{-x_{n}\left|\binom{V_{n+1}}{V_{n}}\right|\left|V_{n}\right|},
$$

implying $x_{n} \leq x_{n+1}$ as desired.
Proof of Theorem 1.16. Assume that $\mathbf{V}$ is a good, homogeneous ssee, and let $\mathcal{P}$ be a hereditary property of $[0,1]$-decorations of $\mathbf{V}$. Set $x_{n}=\frac{\operatorname{ex}\left(V_{n}, \mathcal{P}_{n}\right)}{\left|V_{n}\right|}$.

Suppose first of all that $x_{n} \rightarrow \pi(\mathcal{P})$ as $n \rightarrow \infty$. Thus assumption (a) from Corollary 1.13 is satisfied. We show that assumption (b) is satisfied as well, whence our claimed volume estimate for $\mathcal{P}_{n}$ is immediate. Fix $\varepsilon>0$. We may assume that $\pi(\mathcal{P})>\varepsilon$, else we have nothing to show. Now, by the monotonicity of $\left(x_{n}\right)_{n \in \mathbb{N}}$ established in Proposition 1.15, there exists a constant $N \in \mathbb{N}$ such that
$x_{N}>\pi(\mathcal{P})-\frac{\varepsilon}{3}$. In particular there exists $\delta_{1}=\delta_{1}(N, \varepsilon)>0$ such that if $a \in \operatorname{Box}\left(V_{N}\right)$ satisfies $\left.\operatorname{Ent}(a) \leq\left(\pi(\mathcal{P})-\frac{2 \varepsilon}{3}\right) \right\rvert\, V_{N}$, then we have $\operatorname{vol}\left(a \backslash \mathcal{P}_{N}\right)>\delta_{1}$.

Consider a box $b \in \operatorname{Box}\left(V_{n}\right)$ with $\operatorname{Ent}(b) \leq(\pi(\mathcal{P})-\varepsilon)\left|V_{n}\right|$ for some $n \geq N$. Let $B$ be the family of $\phi \in\binom{V_{n}}{V_{N}}$ such that $\operatorname{Ent}\left(b_{\downarrow \phi}\right) \leq\left(\pi(\mathcal{P})-\frac{2 \varepsilon}{3}\right)\left|V_{N}\right|$. By our observation in the previous paragraph, $\phi \in B$ implies $\operatorname{vol}\left(b_{\downarrow \phi} \backslash \mathcal{P}_{N}\right)>\delta_{1}$. Now, by the homogeneity of $\mathbf{V}$ and the fact the volume of a box is the product of its projections, we have

$$
\begin{aligned}
\left|\binom{V_{n}}{V_{N}}\right|\left|V_{N}\right|(\pi(\mathcal{P})-\varepsilon) & \geq\left|\binom{V_{n}}{V_{N}}\right| \frac{\left|V_{N}\right|}{\left|V_{n}\right|} \operatorname{Ent}(b)=\sum_{\phi \in\binom{V_{n}}{V_{N}}} \operatorname{Ent}\left(b_{\downarrow \phi}\right) \\
& \geq\left|\binom{V_{n}}{V_{N}} \backslash B\right|\left|V_{N}\right|\left(\pi(\mathcal{P})-\frac{2 \varepsilon}{3}\right),
\end{aligned}
$$

implying the existence of a constant $\delta_{2}=\delta_{2}(\varepsilon, \mathcal{P})>0$ such that $|B| \geq \delta_{2}\left|\binom{V_{n}}{V_{N}}\right|$. (Explicitly, $\delta_{2}=\frac{\varepsilon}{3 \pi(\mathcal{P})-2 \varepsilon}$ will do.) It follows that

$$
\mathbb{E}_{\phi} \operatorname{vol}\left(b_{\downarrow \phi} \backslash \mathcal{P}_{N}\right) \geq \delta_{1} \mathbb{P}(\phi \in B) \geq \delta_{1} \delta_{2} .
$$

Setting $\eta=\delta_{1} \delta_{2}$, we have that assumption (b) from Corollary 1.13 is satisfied, and we are done in this case.

Suppose now instead that $x_{n} \rightarrow \infty$ as $n \rightarrow \infty$. For every $C>0$, there exists a constant $N=N(C, \mathcal{P}) \in \mathbb{N}$ such that $x_{N}>2 C+2$. In particular there exists $\delta_{1}=\delta_{1}(C)>0$ such that if $a \in \operatorname{Box}\left(V_{N}\right)$ satisfies $\operatorname{Ent}(a) \leq(2 C+1)\left|V_{N}\right|$, then we have $\operatorname{vol}\left(a \backslash \mathcal{P}_{N}\right)>\delta_{1}$.

Consider a box $b \in \operatorname{Box}\left(V_{n}\right)$ with $\operatorname{Ent}(b) \leq 2 C\left|V_{n}\right|$ for some $n \geq N$. Let $B$ be the family of $\phi \in\binom{V_{n}}{V_{N}}$ such that $\operatorname{Ent}\left(b_{\downarrow \phi}\right) \leq(2 C+1)\left|V_{N}\right|$. By homogeneity of $\mathbf{V}$ and the fact the volume of a box is the product of its projections, we have

$$
\left|\binom{V_{n}}{V_{N}}\right|\left|V_{N}\right| 2 C \geq\left|\binom{V_{n}}{V_{N}}\right| \frac{\left|V_{N}\right|}{\left|V_{n}\right|} \operatorname{Ent}(b)=\sum_{\phi \in\binom{V_{n}}{V_{N}}} \operatorname{Ent}\left(b_{\downarrow}\right) \geq\left|\binom{V_{n}}{V_{N}} \backslash B\right|\left|V_{N}\right|(2 C+1)
$$

implying the existence of a constant $\delta_{2}=\delta_{2}(C)>0$ such that $|B| \geq \delta_{2}\left|\binom{V_{n}}{V_{N}}\right|$ (Explicitly, $\delta_{2}=\frac{1}{2 C+1}$ will do.) Now we have (by our observation in the paragraph above)

$$
\mathbb{E}_{\phi} \operatorname{vol}\left(b_{\downarrow \phi} \backslash \mathcal{P}_{N}\right) \geq \delta_{1} \mathbb{P}(\phi \in B) \geq \delta_{1} \delta_{2} .
$$

Thus we have shown the following: for all $n \geq N$, and all boxes $b \in \operatorname{Box}\left(V_{n}\right), \mathbb{E}_{\phi} \operatorname{vol}\left(b_{\downarrow} \backslash \mathcal{P}_{N}\right)<\delta_{1} \delta_{2}$ implies Ent $(b)>2 C\left|V_{n}\right|(\dagger)$.

We now apply Theorem 1.12 to the good ssee $\mathbf{V}$ and the hereditary property $\mathcal{P}$ with parameters $N \in \mathbb{N}$ and $\varepsilon=\varepsilon(C)=\min \left(C, \delta_{1} \delta_{2}\right)>0$ : there exists $n_{0}=n_{0}(C)>N$ such that for all $n \geq n_{0}$ there exists a collection $\mathcal{C} \subseteq \operatorname{Box}\left(V_{n}\right)$ satisfying properties (i)-(iii) from the statement of Theorem 1.12. By $(\dagger)$ established above and our choice of $\varepsilon>0$, property
(ii) implies that for every $c \in \mathcal{C}$, $\operatorname{Ent}(c)>2 C\left|V_{n}\right|$. Then properties (i), (iii) and our choice of $\varepsilon \leq C$ together yield that for all $n \geq n_{0}(C)$,

$$
\operatorname{vol}\left(\mathcal{P}_{n}\right) \leq \operatorname{vol}\left(\bigcup_{c \in \mathcal{C}} c\right) \leq \sum_{c \in \mathcal{C}} e^{-\operatorname{Ent}(c)} \leq|\mathcal{C}| e^{-2 C\left|V_{n}\right|} \leq e^{(\varepsilon-2 C)\left|V_{n}\right|} \leq e^{-C\left|V_{n}\right|}
$$

Since $C>0$ was arbitrary, it follows that $\operatorname{vol}\left(\mathcal{P}_{n}\right)=e^{-\omega\left(\left|V_{n}\right|\right)}$, as claimed.

## 3 | APPLICATIONS

## 3.1 | Functions from hypercubes into [0, 1]

Let $Q_{n}=\{0,1\}^{n}$ denote the $n$-dimensional hypercube. Consider the sequence of sets $\left(Q_{n}\right)_{n \in \mathbb{N}}$ together with for every $N \leq n$ the collection of injections $\phi: Q_{N} \rightarrow Q_{n}$ obtained by choosing an arbitrary element $\mathbf{u} \in Q_{n}$ and an arbitrary set $S=\left\{s_{1}, s_{2}, \ldots, s_{N}\right\}$ of integers with $1 \leq s_{1}<s_{2}<\cdots<s_{N} \leq n$ and letting

$$
\phi(\mathbf{v})_{i}= \begin{cases}v_{j} & \text { if } i=s_{j} \\ u_{i} & \text { otherwise }\end{cases}
$$

It is an easy exercise to see that this constitutes a good, homogeneous ssee, which we denote by $\mathbf{Q}$.

Remark 3.1. In fact, we still get a good, homogeneous ssee if we consider any of the other natural collections of embeddings $\phi$ on hypercubes, such as for example all the embeddings $\phi$ obtained by selecting $\mathbf{u} \in Q_{n}$ and an injection $\psi:[N] \rightarrow[n]$ and letting

$$
\phi(\mathbf{v})_{i}= \begin{cases}v_{j}+u_{i}(\bmod 2) & \text { if } i=\psi(j), \\ u_{i} & \text { otherwise }\end{cases}
$$

but we do not pursue this here.
Let $c \in \mathbb{R}_{>0}$. Recall that a function $f: Q_{n} \rightarrow[0,1]$ is $c$-Lipschitz if changing a coordinate of $\mathbf{u} \in\{0,1\}^{n}$ changes the values of $f(\mathbf{u})$ by at most $c$. Our aim in this subsection is to show that the probability a random function $f: Q_{n} \rightarrow[0,1]$ is $c$-Lipschitz is $c^{2^{n+o(1)}}$ :

Theorem 3.2. Let $c \in(0,1)$. Let $f: Q_{n} \rightarrow[0,1]$ be a random function chosen according to the uniform measure on $[0,1]^{Q_{n}}$. Then

$$
\mathbb{P}(f \text { is } c \text {-Lipschitz })=c^{2^{n}+o\left(2^{n}\right)} .
$$

We prove Theorem 3.2 via a simple extremal entropy result, Theorem 3.3, from which Theorem 3.2 can be easily deduced via Theorem 1.16. Fix $c \in(0,1)$. Clearly, the collection of functions $Q_{n} \rightarrow[0,1]$ can be identified with the set of $[0,1]$-decorations of $Q_{n}$. Consider the
hereditary property $\mathcal{P}$ of $[0,1]$-decorations of the (good, homogeneous) ssee $\mathbf{Q}$ corresponding to being $c$-Lipschitz.

Theorem 3.3. $\operatorname{ex}\left(Q_{n}, \mathcal{P}_{n}\right)=-\left|Q_{n}\right| \log c$.
Proof. For the upper bound, observe that the box $b=[0, c]^{Q_{n}}$ lies wholly inside $\mathcal{P}_{n}$ and has entropy exactly $-\left|Q_{n}\right| \log (c)$.

For the lower bound, fix $\varepsilon>0$ with $c+\varepsilon<1$. Let $b \in \operatorname{Box}\left(Q_{n}\right)$ be such that $\operatorname{Ent}(b)=\operatorname{ex}\left(Q_{n}, \mathcal{P}_{n}\right)$ and $\operatorname{vol}\left(b \backslash \mathcal{P}_{n}\right)=0$. Then $b=\prod_{\mathbf{x} \in Q_{n}} A_{\mathbf{x}}$, where the $A_{\mathbf{x}}$ are measurable subsets of $[0,1]$. Set $\mathcal{X}=\left\{\mathbf{x}:\left|A_{\mathbf{x}}\right|>c\right\}$. Consider $\mathbf{x} \in \mathcal{X}$, and assume $\left|A_{\mathbf{x}}\right|=\ell$. Then for any $\eta>0$, there exists an interval $I=[u+\eta, u+\ell-\eta] \subseteq[0,1]$ such that both $A_{\mathbf{x}} \backslash[0, u+\ell-\eta]$ and $A_{\mathbf{x}} \backslash[u+\eta, 1]$ have strictly positive measure. Let $\mathbf{y} \in Q_{n}$ be obtained by modifying exactly one coordinate of $\mathbf{x}$. Since $\operatorname{vol}\left(b \backslash \mathcal{P}_{n}\right)=0$ and $A_{\mathbf{x}} \backslash[0, u+\ell-\eta]$ has a positive measure, the definition of $c$-Lipschitz implies that $A_{\mathbf{y}} \cap[0, u+\ell-\eta-c]$ must have zero measure. Similarly $A_{\mathbf{x}} \backslash[u+\eta, 1]$ having a positive measure implies $A_{\mathbf{y}} \cap[u+\eta+c, 1]$ has zero measure. Overall we obtain $\left|A_{\mathbf{y}} \cap[u+\ell-\eta-c, u+\eta+c]\right|=\left|A_{\mathbf{y}}\right|$. Since $\eta>0$ was arbitrarily chosen, this implies in fact that up to a zero-measure set $A_{\mathbf{y}}$ is contained inside the interval $[u+\ell-c, u+c]$. In particular we must have $\ell \leq 2 c$, and $\left|A_{\mathbf{y}}\right| \leq(2 c-\ell)<c$.

Now partition $Q_{n}$ into pairs $\{\mathbf{v} \times\{0\}, \mathbf{v} \times\{1\}\}$, with $\mathbf{v}$ running over all possible choices $\mathbf{v} \in Q_{n-1}$. By the above, we have that each such pair $\{\mathbf{x}, \mathbf{y}\}$ contains at most one element of $\mathcal{X}$. Moreover if this element is $\mathbf{x}$ and satisfies $\left|A_{\mathbf{x}}\right|=\ell$, then $\left|A_{\mathbf{y}}\right| \leq(2 c-\ell)$, and

$$
\left|A_{\mathbf{x}}\right| \cdot\left|A_{\mathbf{y}}\right| \leq \ell \cdot(2 c-\ell)<c^{2}
$$

Thus we have

$$
\operatorname{vol}(b)=\prod_{\mathbf{v} \in Q_{n-1}}\left|A_{\mathbf{v} \times\{0\}}\right|\left|A_{\mathbf{v} \times\{1\}}\right| \leq c^{2^{n}-2|\mathcal{X}|} c^{2|X|}=c^{2^{n}},
$$

with equality if and only if $\mathcal{X}=\varnothing$ and for every $\mathbf{x} \in Q_{n}$ we have $\left|A_{\mathbf{x}}\right|=c$. This shows that $\operatorname{ex}\left(Q_{n}, \mathcal{P}_{n}\right) \geq-\left|Q_{n}\right| \log c$, as claimed.

Proof of Theorem 3.2. Let $\mathcal{P}$ denote, as above, the property of being $c$-Lipschitz, viewed as a hereditary property of $[0,1]$-decorations of the good homogeneous ssee $\mathbf{Q}$. By Theorem 3.3, $\pi(\mathcal{P})=-\log c$, whence the volume estimate

$$
\operatorname{vol}\left(\mathcal{P}_{n}\right)=e^{(\log (c)+o(1))\left|Q_{n}\right|}=c^{2^{n}+o\left(2^{n}\right)}
$$

follows immediately from Theorem 1.16, implying the desired estimate for $\mathbb{P}\left(f \in \mathcal{P}_{n}\right)$.

## 3.2 | Metric polytopes

Given a set $S$ and a positive integer $s$, we write $S^{(s)}$ for the collection of subsets of $S$ of size $s$. Let $K_{n}=(V, E)$ denote the complete graph on the vertex-set $V=V\left(K_{n}\right)=[n]$ with edge-set $E=E\left(K_{n}\right)=[n]^{(2)}$. It is an easy exercise to check that the sequence $\left(E\left(K_{n}\right)_{n \in \mathbb{N}}\right)$ together with
the collection of embeddings $\phi$ corresponding to graph isomorphisms $K_{N} \rightarrow K_{n}$ constitutes a good homogeneous ssee, which we denote by $\mathbf{K}$.

A $[0,1]$-decoration $d \in[0,1]^{E\left(K_{n}\right)}$ of the edges of $K_{n}$ can be seen as an assignment of distances to pairs of vertices of $K_{n}$. Let $\mathcal{M}_{n}$ denote the collection of such $d$ which satisfy the triangle inequality (and for which $\left(V\left(K_{n}\right), d\right.$ ) is thus a metric space). The body $\mathcal{M}_{n}$ is known as the metric polytope. The property $\mathcal{M}=\left(\mathcal{M}_{n}\right)_{n \in \mathbb{N}}$ is hereditary, since a subset of a metric space is itself a metric space. Thus we can use an easy extremal argument together with our main results to estimate the volume of $\mathcal{M}_{n}$, and thereby prove a weak form of a recent theorem of Kozma, Meyerovitch, Peled and Samotij [24].

Theorem 3.4 (Rough estimate for the volume of the metric polytope). $\operatorname{vol}\left(\mathcal{M}_{n}\right)=\left(\frac{1}{2}\right)^{\binom{n}{2}+o\left(n^{2}\right)}$.
Theorem 3.4 will follow from the following extremal result:
Theorem 3.5. For all $n \geq 3, \operatorname{ex}(\mathbf{K}, \mathcal{M})_{n}=\binom{n}{2} \log 2$.

Proof. For the upper bound, observe that the box $b_{n}=\left[\frac{1}{2}, 1\right]^{E\left(K_{n}\right)} \subseteq \mathcal{M}_{n}$ and has entropy $\binom{n}{2} \log 2$ (since clearly for every $d \in b_{n}$, the associated assignment of distances to the edges of $K_{n}$ satisfies the triangle inequality). Thus $\operatorname{ex}(\mathbf{K}, \mathcal{M})_{n} \leq\binom{ n}{2} \log 2$ for all $n \geq 3$.

For the lower bound, consider a simple box $b \in \operatorname{Box}\left(E\left(K_{3}\right)\right)$ with $\operatorname{vol}\left(b \backslash \mathcal{M}_{3}\right)=0$. Then $b=I_{1} \times I_{2} \times I_{3}$, where $I_{i}$ is a union of nonempty intervals (and corresponds to the edge $[3] \backslash\{i\}$ of $K_{3}$ ). Set $a_{i}=\min I_{i}$ and $b_{i}=\max I_{i}$. Clearly, we have $b \subseteq \prod_{i}\left[a_{i}, b_{i}\right] \subseteq \prod_{i}\left[a_{i}, 1\right]$. In particular we have

$$
\begin{equation*}
\operatorname{vol}(b) \leq \prod_{i}\left(b_{i}-a_{i}\right) \quad \text { and } \quad \operatorname{vol}(b) \leq \prod_{i}\left(1-a_{i}\right) \tag{6}
\end{equation*}
$$

Let $A=\sum_{i} a_{i}$. Since $\operatorname{vol}\left(b \backslash \mathcal{M}_{3}\right)=0$, it follows by the triangle inequality that for all $i \in[3], b_{i} \leq a_{i+1}+a_{i+2}$ (where the indices are taken modulo 3). Summing over all $i$ and subtracting $A$ from both sides, we get $\sum_{i}\left(b_{i}-a_{i}\right) \leq A$. Simple calculus then tells us that for $A$ fixed, we have $\prod_{i}\left(b_{i}-a_{i}\right) \leq(A / 3)^{3}$. On the other hand, again by simple calculus, for $A$ fixed we have $\prod_{i}\left(1-a_{i}\right) \leq(1-A / 3)^{3}$. Combining these two bounds with (6), we get that

$$
\operatorname{vol}(b) \leq \min \left((A / 3)^{3},(1-A / 3)^{3}\right)=2^{-3}
$$

As $b$ was an arbitrary simple box with $\operatorname{vol}\left(b \backslash \mathcal{M}_{3}\right)=0$, this implies $\operatorname{ex}(\mathbf{K}, \mathcal{M})_{3}=\binom{3}{2} \log 2$. Since $\mathbf{K}$ is homogeneous, it follows from Proposition 1.15 that $\operatorname{ex}(\mathbf{K}, \mathcal{M})_{n} \geq\binom{ n}{2} \log 2$ for all $n \geq 3$, as required.

Proof of Theorem 3.4. Immediate from Theorems 3.5 and 1.16 applied to the good homogeneous ssee K.

The problem of estimating $\operatorname{vol}\left(\mathcal{M}_{n}\right)$ has been previously considered by several other researchers, who obtained significantly stronger estimates than Theorem 3.4. As observed by Kozma, Meyerovitch, Peled and Samotij, the problem of estimating $\operatorname{vol}\left(\mathcal{M}_{n}\right)$ is related to the problem of estimating the number of metric spaces on $n$ points with integer distances, which was studied by Mubayi and Terry [30] using the container method. Balogh and Wagner [7, Theorem 3.7] used the container method to show $\operatorname{vol}\left(\mathcal{M}_{n}\right) \leq\left(\frac{1}{2}\right)^{\binom{n}{2}+n^{11 / 6+o(1)}}$. Finally, in a recent and impressive paper using entropy techniques, Kozma, Meyerovitch, Peled and Samotij obtained much more precise estimates on $\operatorname{vol}\left(\mathcal{M}_{n}\right)$ : they proved in [24, Theorem 1.2] that

$$
\left(\frac{1}{2}\right)^{\binom{n}{2}} e^{\frac{n^{3 / 2}}{6}+o\left(n^{3 / 2}\right)} \leq \operatorname{vol}\left(\mathcal{M}_{n}\right) \leq\left(\frac{1}{2}\right)^{\binom{n}{2}} e^{O\left(n^{3 / 2}\right)} .
$$

They also showed in Section 5.3 of the paper, in joint work with Morris, how using the more precise container theorems of [5] (rather than the simple containers of [34] that underpin Theorem 2.1 and all the results in this paper) could be made to yield slightly weaker upper bounds of $\left(\frac{1}{2}\right)^{\binom{n}{2}} e^{O\left(n^{3 / 2}(\log n)\right)}$ on $\operatorname{vol}\left(\mathcal{M}_{n}\right)$, improving on the earlier container results of Balogh and Wagner. The point of Theorem 3.4 is thus not to prove a new or optimal upper bound, but rather to illustrate how the results of this paper give a simple, streamlined approach to such problems.

## 3.3 | Weighted graphs

A [0, 1]-decoration $w$ of the edges of $K_{n}$ may be identified may be viewed as an assignment of weights $w(e) \in[0,1]$ to the edges $e \in E\left(K_{n}\right)$. Given such a decoration, we may define the weight of a set $S \subseteq V\left(K_{n}\right)$ as $w(S):=\sum_{e \in S^{(2)}} w(e)$. For a fixed integer $s \geq 2$ and a real number $r \in\left(0,\binom{s}{r}\right]$, let $\mathcal{P}(s, r)$ be the hereditary property of $[0,1]$-decorations of the edges of $\mathbf{K}$ corresponding to having no $s$-set of vertices whose weight exceeds $r$.

Theorem 3.6. For all $n \geq s$ we have $\operatorname{ex}\left(E\left(K_{n}\right), \mathcal{P}_{n}(s, r)\right)=\binom{n}{2} \log \left(\frac{\binom{s}{2}}{r}\right)$, with equality uniquely attained up to zero-measure sets by the box $\left[0, \frac{r}{\binom{s}{2}}\right]^{E\left(K_{n}\right)}$.

Proof. For the upper bound, note that $\left[0, \frac{r}{\binom{s}{2}}\right]^{E\left(K_{n}\right)}$ is a box lying wholly inside $\mathcal{P}_{n}$ and having the claimed volume. For the lower bound, note first of all that up to a zeromeasure set, any entropy-maximiser $b$ for $\mathcal{P}_{n}$ must be of the form $b=\prod_{e \in E\left(K_{n}\right)}[0, w(e)]$, where $w: E\left(K_{n}\right) \rightarrow[0,1]$ is an edge-weighting from $\mathcal{P}_{n}$. By averaging the weight $w(S)$ over all $s$-sets $S$ and using the fact that $w \in \mathcal{P}(s, r)$, we see that

$$
\binom{n-2}{s-2} w\left(V\left(K_{n}\right)\right)=\sum_{S \in[n]^{(s)}} w(S) \leq r\binom{n}{s} .
$$

In particular, $\sum_{e \in E\left(K_{n}\right)} w(e)=w\left(K_{n}\right) \leq \frac{r}{\binom{s}{2}}\binom{n}{2}$. It then follows from the AM-GM inequality that

$$
\operatorname{vol}(b)=\prod_{e \in E\left(K_{n}\right)} w(e) \leq\left(\frac{r}{\binom{s}{2}}\right)^{\binom{n}{2}}
$$

which gives the required lower bound on the entropy of $b$. Furthermore the AM-GM inequality also implies that equality is attained if and only if $w(e)=\left[0, r /\binom{s}{2}\right]$ for every $e \in E\left(K_{n}\right)$, that is, if and only if up to a zero-measure set $b$ is equal to the box $[0, w]^{E\left(K_{n}\right)}$, as claimed.

Corollary 3.7. $\operatorname{vol}\left(\mathcal{P}_{n}(s, r)\right)=\left(\frac{r}{\binom{s}{2}}+o(1)\right)^{\binom{n}{2}}$.

Proof. Theorem 3.6 shows $\pi(\mathcal{P}(s, r))=r /\binom{s}{2}$. The result is then immediate from an application of Theorem 1.16 to the good, homogeneous ssee $\mathbf{K}$.

We should note here that Mubayi and Terry [31, 32] considered the related problem of maximising the product of edge-multiplicities in multigraphs in which every $s$-set of vertices spans at most $r$ edges. In this case the fact that edge-multiplicities have to be positive integers completely changes the nature of the problem, which becomes highly nontrivial (see [13, 15] for recent progress on the Mubayi-Terry problem).

## 4 | ENTROPY OF [k]-DECORATED GRAPH LIMITS

As another application of our work, we prove a generalisation of the result of Hatami, Janson and Szegedy on the entropy of graph limits to [ $k$ ]-decorated graph limits. Hatami, Janson and Szegedy defined and studied the entropy of a graphon in [22]. They used this notion to give an alternative proof of the Alekseev-Bollobás-Thomason theorem [2, 9] and to describe the typical structure of a graph in a hereditary property. The Hatami-Janson-Szegedy notion of entropy can be viewed as a graphon analogue of the classical notion of the entropy of a discrete random variable. Generalising their result to [ $k$ ]-decorated graphons was one of the original motivations for our foray into container theory (as containers allow for an easy transfer of certain results to the limit setting). In fact, we sought unsuccessfully to obtain a generalisation to [0, 1]-decorated graph limits, which we now define.

Let $\mathcal{K}$ be a compact second-countable Hausdorff space. A $\mathcal{K}$-decorated graph is a function $w: E\left(K_{n}\right) \rightarrow \mathcal{K}$ assigning to each edge of the complete graph $K_{n}$ a label from $\mathcal{K}$. In [28], Lovász
and Szegedy initiated the study of the limits of sequences of $\mathcal{K}$-decorated graphs. Generalising the well-established theory of graph limits, given a $\mathcal{K}$-decorated graph $G$ they defined homomorphism densities $t(F, G)$ for $C[\mathcal{K}]$-decorated graphs $F$ in $G$, where $C[\mathcal{K}]$ is the collection of all continuous functions $\mathcal{K} \rightarrow \mathbb{R}$. They defined convergence relative to this notion of homomorphism density, and showed the limit objects in this theory were $\mathcal{K}$-graphons, which are symmetric measurable functions $W:[0,1]^{2} \rightarrow \mathcal{M}(\mathcal{K})$, where $\mathcal{M}[\mathcal{K}]$ denotes the set of Borel probability measures on $\mathcal{K}$. Further contributions to the study of $\mathcal{K}$-decorated graph limits were made by Kunszenti-Kovács, Lovász and Szegedy [25], who defined a modified notion of cut distance (called jumble distance, which they used to provide weak regularity lemma and a counting lemma for $\mathcal{K}$-graphons) and showed that the space of $\mathcal{K}$-decorated graph limits was closed under the homomorphism density notion of convergence. However a number of questions, such as the uniqueness of the representation of the limits of $\mathcal{K}$-decorated graph sequences or the compactness of that space under their modified notion of cut distance, remain open.

Our goal was to extend the results of Hatami-Janson-Szegedy on graphon entropy to [0, 1]decorated graphons. However, we were unable to show that every sequence of [0, 1]-decorated graphs contains a convergent subsequence of [0, 1]-decorated graphs. Thus we had to content ourselves with proving a generalisation of Hatami-Janson-Szegedy to the easier setting of [k]-decorated graphs, which we now give. Before we can state our results, we must recall the definitions of templates and realisations from Section 2, and also recall from [18] that the entropy of a $[k]$-colouring template $t$ for $K_{n}$ is

$$
\operatorname{Ent}(t):=\log _{k} \prod_{e \in E\left(K_{n}\right)}|t(e)|
$$

Note that for any template $t$ we have $0 \leq \operatorname{Ent}(t) \leq\binom{ n}{2}$, and the number of realisations of $t$ is exactly $|\langle t\rangle|=k^{\operatorname{Ent}(t)}$. For a hereditary property $\mathcal{P}$ of $[k]$-decorated graphs, we defined

$$
\operatorname{ex}\left(n, \mathcal{P}_{n}\right):=\max \left\{\operatorname{Ent}(t): t \text { is a } k \text {-colouring template for } K_{n} \text { with }\langle t\rangle \subseteq \mathcal{P}_{n}\right\}
$$

Finally, let us recall the [ $k$ ]-decorated graph analogue of Theorem 1.16 from [18].
Theorem 4.1 (Falgas-Ravry et al. [18, Corollary 2.15]). Let $\mathcal{P}$ be a hereditary property of $[k]$-decorated graphs with $\mathcal{P}_{n} \neq \varnothing$ for every $n \in \mathbb{N}$ and let $\varepsilon>0$ be fixed. There exists $n_{0} \in \mathbb{N}$ such that for all $n \geq n_{0}$,

Given a [k]-graphon $W$ and $i \in[k]$, we denote by $W_{i}(x, y):=W(x, y)(i)$ the probability of $\{i\}$ under the probability measure $W(x, y)$ on $[k]$. As noted by Lovász and Szegedy [28, Example 2.8] each $W_{i}$ is a graphon. Given a probability measure $P$ on $[k]$ with $P(i)=p_{i}$, we define the $k$-ary entropy of $P$ to be

$$
h_{k}(P):=\sum_{i \in[k]}-p_{i} \log _{k} p_{i}
$$

Then the entropy of a $[k]$-graphon $W$ is

$$
\operatorname{Ent}(W):=\iint_{[0,1]^{2}} h_{k}\left(W_{k}(x, y)\right) d A
$$

Note that $0 \leq \operatorname{Ent}(W) \leq 1$. For $k=2$ our definition of decorated graphon entropy coincides with that of Hatami, Janson and Szegedy. Given a property $\mathcal{P}$ of $[k]$-decorated graphs, we denote by $\hat{\mathcal{P}}$ the closure under the cut norm (see Section 4.1 for a definition of the cut norm) of the collection of [k]-graphons that can be obtained as a limit of a convergent sequence of elements of $\mathcal{P}$. We can at last state the main result of this section.

Theorem 4.2. Let $\mathcal{P}$ be a hereditary property of $[k]$-decorated graphs and let $m(\hat{\mathcal{P}}):=\max _{W \in \hat{\mathcal{P}}} \operatorname{Ent}(W)$. Then

$$
\lim _{n \rightarrow \infty} \frac{\log _{k}\left|\mathcal{P}_{n}\right|}{\binom{n}{2}}=m(\hat{\mathcal{P}}) .
$$

Given Theorem 4.1, the theorem above is equivalent to the assertion that $\pi(\mathcal{P})=m(\hat{\mathcal{P}})$, which in fact is what we shall prove.

## 4.1 | A cut distance for [k]-graphons

We require convergence in our proof, but rather than using convergence with respect to homomorphisms, we use convergence with respect to an appropriately defined cut distance. Frieze and Kannan [19] introduced a cut norm $\|\cdot\|_{\square}$ that has become central to the theory of graph limits (see [23, Section 4] for an overview of the history of the cut norm in other contexts). The cut norm of a graphon $W$ is

$$
\|W\|_{\square}:=\sup _{S, T \subseteq[0,1]}\left|\int_{S \times T} W(x, y) d x d y\right|,
$$

where the supremum is over all pairs $(S, T)$ of measurable subsets of $[0,1]$. If $U$ and $W$ are graphons, then

$$
d_{\square}(U, W):=\|U-W\|_{\square}=\sup _{S, T \subseteq[0,1]}\left|\int_{S \times T}(U(x, y)-W(x, y)) d x d y\right| .
$$

Given a measure-preserving transformation $\varphi:[0,1] \rightarrow[0,1]$, we define $W^{\varphi}$ by $W^{\varphi}(x, y):=W(\varphi(x), \varphi(y))$. The cut distance between $U$ and $W$ is

$$
\delta_{\square}(U, W):=\inf _{\varphi:[0,1] \rightarrow[0,1]} d_{\square}\left(U, W^{\varphi}\right),
$$

where the infimum is taken over all measure-preserving transformations $\varphi:[0,1] \rightarrow[0,1]$. We introduce an appropriate generalisation of the cut distance here (this was also previously
considered in [25]). If $G$ and $H$ are two [ $k$ ]-decorated graphs with edge labellings $g, h: E\left(K_{n}\right) \rightarrow[k]$, respectively, and with vertex-set [ $n$ ], we define

$$
d_{\square k}(G, H):=\max _{S, T \subseteq[n]} \frac{1}{n^{2}} \sum_{i=1}^{k}\left|\sum_{(u, v) \in S \times T}(\mathbb{1}(g(u v)=i)-1(h(u v)=i))\right| .
$$

If $U$ and $W$ are $[k]$ graphons, we define

$$
d_{\square k}(U, W):=\sup _{S, T \subseteq[0,1]} \sum_{i=1}^{k} \mid \int_{S \times T}\left(U_{i}(x, y)-W_{i}(x, y) d x d y \mid .\right.
$$

We define the cut distance $\delta_{\square k}$ for $[k]$-graphons analogously to the definition for graphons, mutatis mutandis. Letting $\mathcal{W}_{k}$ denote the set of all $[k]$-graphons, we let $\widetilde{\mathcal{W}_{k}}$ denote the quotient of $\mathcal{W}_{k}$ obtained by identifying $U$ and $W$ whenever $\delta_{\square_{k}}(U, W)=0$.

Theorem 4.3. The space $\left(\widetilde{\mathcal{W}_{k}}, \delta_{\square k}\right)$ is compact.
The proof is essentially identical to the proof of compactness with respect to cut distance for ordinary graphons as given by the original argument of Lovász and Szegedy [27, Theorem 5.1]. Therefore we only give a brief sketch here.

Sketch proof of Theorem 4.3. First note that a weak regularity lemma for $[k]$ graphons follows very quickly from the weak regularity lemma for ordinary graphons (Lemma 3.1 in [27]), simply by running it for each $W_{i}(x, y), i \in[k]$ simultaneously. Now given a sequence of $[k]$-graphons $W_{1}, W_{2}, \ldots \in \widetilde{\mathcal{W}_{k}}$, using the weak regularity lemma, we find for every $n \in \mathbb{N}$ a sequence of step-functions $W_{n, \ell}, \ell \in \mathbb{N}$, converging to $W_{n}$. Now for each $\ell$ we find a subsequence $n_{i}, i \in \mathbb{N}$ for which $W_{n_{i} \ell}$ converges in cut distance to a $[k]$ graphon $U_{\ell}$. Then by the Martingale Convergence Theorem, the sequence $\left(U_{\ell}\right)$ converges to a limit $U$. Finally one can show $\delta_{\square_{k}}\left(W_{i}, U\right) \rightarrow 0$ using a $3 \varepsilon$-argument.

## 4.2 | Going between templates and graphons

Fundamental to the theory of graph limits is a natural way of obtaining a graphon from a given graph, and conversely (via sampling) a way of obtaining a graph on $n$ vertices from a given graphon. These transformations respect homomorphism densities and cut distance, and in particular, with probability one, a sequence of $n$-vertex graphs $\left(G_{n}\right)_{n \in \mathbb{N}}$ sampled from $W$ converges to $W$ itself (as established in [10]). Similarly here, we obtain a way of going between the discrete and limit objects. The only property we require is that this transformation respects entropy, which as we will see follows easily from the definitions.

Given a set of $n$ points $x_{1}, \ldots, x_{n}$ from $[0,1]$ and a $[k]$-graphon $W$, we may define a $k$-colouring template for $K_{n}, t_{W}\left[x_{1}, \ldots, x_{n}\right]$, by setting $t(i j):=\left\{c \in[k]: \mathbb{P}\left(W\left(x_{i}, x_{j}\right)=c\right)>0\right\}$. Further we may define a random $k$-colouring $c_{W}\left[x_{1}, \ldots, x_{n}\right]$ by setting $c(i j)$ to be a random
colour from $[k]$ drawn according to the probability distribution given by $W\left(x_{i}, x_{j}\right)$. We define the $W$-random template $t_{W}(n)$ and the $W$-random colouring $c_{W}(n)$ by selecting $x_{1}, \ldots, x_{n}$ uniformly at random from $[0,1]$ and then taking the resulting (induced) $k$-colouring template and random $k$-colouring, respectively.

Our $W$-random templates and colourings give us a way of going from [k]-graphons to $k$ colouring templates and $k$-colourings of $E\left(K_{n}\right)$. We can also go in the other direction: first divide $[0,1)$ into intervals $I_{i}:=[(i-1) / n, i / n)$ for $1 \leq i \leq n$. Given a $k$-colouring template $t$ of $K_{n}$, we may define $W_{t}$ by defining for every $(x, y) \in[0,1]^{2}$ and $c \in[k]$,

$$
\left(W_{t}\right)_{c}:= \begin{cases}\frac{1}{|t(i j)|} 1(c \in t(i j)) & \text { if }(x, y) \in I_{i} \times I_{j}, \quad 1 \leq i, j \leq n, i \neq j \\ \frac{1}{k} & \text { if }(x, y) \in I_{i} \times I_{i}, \quad 1 \leq i \leq n\end{cases}
$$

In other words, for each tile $I_{i} \times I_{j}$ we distribute the mass evenly over the colours which appear in $t(i j)$, and give the uniform distribution to the diagonal tiles $I_{i} \times I_{i}$.

By viewing a $k$-colouring $c$ of $E\left(K_{n}\right)$ as a (zero entropy) template, we may in the same way obtain from it a $[k]$-graphon $W_{c}$. Thus we may go in a natural way from properties of colourings to properties of decorated graphons, and vice versa.

Note that for all $k$ and $n$, and every $k$-colouring template $t$ for $K_{n}$, we have $\binom{n}{2} \operatorname{Ent}\left(W_{t}\right)=\operatorname{Ent}(t)+\frac{n-1}{2}$. In particular for a template $t$ and its associated $[k]$-decorated graphon $W_{t}$ we have

$$
\begin{equation*}
\operatorname{Ent}\left(W_{t}\right)=\frac{\operatorname{Ent}(t)+O(n)}{\binom{n}{2}} \tag{7}
\end{equation*}
$$

Furthermore in the reverse direction, given a $[k]$-decorated graphon $W,\binom{n}{2} \operatorname{Ent}(W)-\frac{n-1}{2}$ is exactly the expected value of the discrete $k$-ary entropy of the $W$-random colouring model $t_{W}(n)$.

## 4.3 | Proof of main result

Proof of Theorem 4.2. For each $n \in \mathbb{N}$ take an extremal template $t_{n}$ which maximises $\operatorname{ex}\left(n, \mathcal{P}_{n}\right)$. We have

$$
\pi(P)=\lim _{n \rightarrow \infty} \frac{\operatorname{ex}\left(n, \mathcal{P}_{n}\right)}{\binom{n}{2}}=\lim _{n \rightarrow \infty} \frac{\operatorname{Ent}\left(t_{n}\right)}{\binom{n}{2}}
$$

Letting $W_{t_{n}}$ be the $[k]$-graphon corresponding to $t_{n}$, we have $\operatorname{Ent}\left(W_{t_{n}}\right)=\left(\operatorname{Ent}\left(t_{n}\right)+O(n)\right) /\binom{n}{2}$ by (7). By Theorem 4.3, there exists a subsequence $\left(W^{j}\right)_{j \in \mathbb{N}}$ of $\left(W_{t_{n}}\right)_{n \in \mathbb{N}}$ which converges to a limit [k]-graphon $W$, where we have $\delta_{\square_{k}}\left(W^{j}, W\right) \rightarrow 0$ as $j \rightarrow \infty$. Now since $\left\langle t_{n}\right\rangle \subseteq \mathcal{P}_{n}$, for any $N$ fixed, the probability that the $W_{t_{n}}$-random colouring of $K_{N}$ is in $\mathcal{P}_{N}$ is $1-O(N / n)=1-o(1)$. It follows that a
$W$-random colouring lies in $\mathcal{P}$ with probability 1 , and thus $W \in \hat{\mathcal{P}}$. Since entropy is a linear functional we thus have $m(\hat{\mathcal{P}}) \geq \operatorname{Ent}(W)=\pi(P)+o(1)$.

Conversely, let $W$ be an entropy maximiser in $\hat{\mathcal{P}}$ with $\operatorname{Ent}(W)=m(\hat{\mathcal{P}})$. For every $n \in \mathbb{N}$, by the linearity of expectation if $n$ points $x_{1}, \ldots, x_{n}$ are chosen uniformly at random from $[0,1]$ then with strictly positive probability

$$
\operatorname{Ent}\left(t_{W}\left[x_{1}, \ldots, x_{n}\right]\right) \geq m(\hat{\mathcal{P}})\binom{n}{2}
$$

Furthermore, as $W \in \hat{\mathcal{P}}$, almost surely $\left\langle t_{W}\left[x_{1}, \ldots, x_{n}\right]\right\rangle \subseteq \mathcal{P}_{n}$, which implies that $\left|\mathcal{P}_{n}\right| \geq k^{m(\hat{\mathcal{P}})\binom{n}{2}}$ for every $n \in \mathbb{N}$. By Theorem 4.1, we have $\left|\mathcal{P}_{n}\right| \leq k^{(\pi(\mathcal{P})+o(1))\binom{n}{2} \text {, and thus }}$ $\pi(\mathcal{P})+o(1) \geq m(\hat{\mathcal{P}})$ as required.

## 5 | CONCLUDING REMARKS

In this paper, we have explored some consequences of the simple versions of the container theorems of Balogh-Morris-Samotij and Saxton-Thomason for the problem of estimating volume or approximating by boxes for certain hereditary bodies. Many problems remain open however.

## 5.1 | Alternative approaches to containers?

Using the Saxton-Thomason simple container theorem as a black box (which, as we stated at the beginning of Section 2, is the result behind Theorem 2.1 and thus the main tool behind all our results), we showed in Theorem 1.12 that hereditary properties of [0, 1]-decorated ssee-s can be approximated by a "small" union of boxes.

A natural question to ask is whether one can go in the other direction: is it possible to obtain a container theorem from purely geometric considerations on approximations of hereditary bodies by boxes? A simplest version of this question is the following: suppose we have a sequence $\left(b_{n}\right)_{n \in \mathbb{N}}$ of bodies with $b_{n} \subseteq[0,1]^{n}$ and every strict projection of $b_{n}$ into $[0,1]^{N}$ (where we use strict projection in the sense of Definition 1.7) is a subset of $b_{N}$. Does it follow (by measure-theoretic/geometric arguments) that for all $n$ sufficiently large there exists a "fine" approximation of $b_{n}$ by a "small" collection of simple boxes?

## 5.2 | Questions about graph limits

In a different direction, we have tried to connect some container-derived results with questions about limit objects. A natural question is, again, whether one can go in the other direction, and derive some finitary container theorems from infinitary arguments about limit objects?

The simplest example of this is perhaps the following: suppose we have a hereditary property $\mathcal{P}$ of $\{0,1\}$-decorations of $E\left(K_{n}\right)$ (i.e., of ordinary graphs). Let $\hat{\mathcal{P}}$ denote the closure of
the family of limits of sequences of graphs from $\mathcal{P}$ under the cut norm. Let $\mathcal{Q}$ denote the collection of graphons that lie at graph distance at most $\varepsilon$ from $\hat{\mathcal{P}}$-this is a closed and hence compact set. Introduce a partial order on $\mathcal{Q}$ by setting $W_{1} \succ W_{2}$ if almost everywhere either $\operatorname{Ent}\left(W_{1}(x, y)\right)>0$ or $W_{1}(x, y)=W_{2}(x, y)$ holds. Then for each $W \in \mathcal{Q}$, let $B(W)$ denote the interior of the collection of $W^{\prime} \in \mathcal{P}$ with $W \succ W^{\prime}$. Clearly the $B(W)$ are open sets in the closed, compact set $\mathcal{P}$. Thus if one could show that they also cover $\mathcal{P}$ it would follow by compactness that there exists some finite set $S$ (with size depending on $\varepsilon$ ) of elements of $\mathcal{Q}$ such that $\bigcup_{W \in S} B(W)=\hat{\mathcal{P}}$. One could then plausibly extract from the graphons in $S$ a small family of containers for $\mathcal{P}_{n}$. This or other approaches to the construction of containers "from the limit" and from purely analytic considerations strike us as an intriguing problem.

With regard to limit objects, the other obvious question is generalising Theorem 4.2 to [ 0,1 ]-decorated graphons. Here our problem is that we did not prove the compactness of the limit space under the cut distance (i.e., we do not have a $[0,1]$-decorated version of Theorem 4.3), and so given a sequence of boxes from $\mathcal{P}_{n}$ we could not extract a subsequence converging to an element of $\hat{\mathcal{P}}$, which we needed to bound $m(\hat{P})$ below. Addressing this issue would immediately extend our results to [ 0,1 ]-decorated graphons and in addition would advance the project of Lovász and Szegedy of building a theory for $\mathcal{K}$-decorated graph limits for second-countable compact Hausdorff spaces $\mathcal{K}$, a worthwhile goal in itself.

## 5.3 | Quality of the container approximation

Can one improve assumption (ii) in Theorem 1.11 (and hence Theorem 1.12)? For instance, could we guarantee that, say

$$
\operatorname{vol}\left(\bigcup_{c \in \mathcal{C}} c\right) \leq C_{\varepsilon} \operatorname{vol}\left(\mathcal{P}_{n}\right)
$$

for some $n$-independent constant $C_{\varepsilon}>1$ ? Or could one show a weaker bound of the form

$$
\operatorname{vol}\left(c \backslash \mathcal{P}_{n}\right)<\varepsilon \operatorname{vol}(c) ?
$$

Putting it in slightly different terms: How fine can we make our approximation of a hereditary body $b$ by simple boxes? There should be a trade-off between the fineness of our approximation and the number of boxes it contains. Is it the case that, for example, the worstcase product of the approximation ratio and the size of the approximation family is bounded below by some function of $\left|V_{n}\right|$ ? Further, what do the bodies that are hardest to approximate look like?

## 5.4 | Relaxing homogeneity

In Theorem 1.16 we obtained a rather clean statement concerning the volume of hereditary properties for homogeneous ssee-s. Homogeneity is a strong condition, however, and it is natural to ask whether it can be relaxed. Explicitly, call an ssee $\mathbf{V}$ almost homogeneous if there
exist constants $C>c>0$ such that for every $n \geq N$, every $x \in V_{n}$ is contained in at least $c\left|\binom{V_{n}}{V_{N}}\right| \frac{\left|V_{N}\right|}{\left|V_{n}\right|}$ and at most $C\left|\binom{V_{n}}{V_{N}}\right| \frac{\left|V_{N}\right|}{\left|V_{n}\right|}$ embeddings $\phi\left(V_{N}\right)$ with $\phi \in\binom{V_{n}}{V_{N}}$.

Can one obtain a version of Theorem 1.16 in which the homogeneity assumption is relaxed to almost homogeneity? This would increase the generality of the results in this paper and allow us to cover some important cases, such as that of the ssee $\mathfrak{I}$ where $\mathfrak{I}_{n}=[n]$ and the embeddings $\phi: \Im_{N} \rightarrow \Im_{n}$ consist of the injections from [ $N$ ] into arithmetic progressions of length $N$ in $[n]$. Another example would be that of permutations, see Section 5.5.

## 5.5 | Containers for thin bodies

In this paper, we have been content with a simple container bound $|\mathcal{C}| \leq e^{\varepsilon\left|V_{n}\right|}$ on the size of the container family $\mathcal{C}$. This is sufficient to estimate the volume of $\mathcal{P}_{n}$ when the maximum volume of a box contained in $\mathcal{P}_{n}$ (up to a zero-measure set) is of order $e^{-\theta\left(\left|V_{n}\right|\right)}$. However, for "thinner" bodies when this extremal volume is of order $e^{-\omega\left(\left|V_{n}\right|\right)}$, our results (more specifically Theorem 1.16) say nothing more precise than $\operatorname{vol}\left(\mathcal{P}_{n}\right)=e^{-\omega\left(I V_{n} \mid\right)}$. This is definitely a limitation of our work-the original container theorems of Balogh-Morris-Samotij and Saxton-Thomason can give much better estimates, but require information on the degree measure (something which, as Saxton and Thomason [34] observe is unnecessary in the case of "thicker" bodies).

There are a number of interesting examples within our framework where more precise estimates would be advantageous-for instance, that of permutations, which we discuss below.

Denote by $S_{n}$ the collection of all permutations of [ $n$ ]. Given a [ 0,1$]$-decoration of [ $n$ ], we may define a permutation in $S_{n}$ as follows. Let $\mathcal{B}$ denote the collection of $x \in[0,1]^{[n]}$ which have at least two coordinates equal. Note that this is a zero-measure set. Given $x \in[0,1]^{[n]} \backslash \mathcal{B}$, we equip $[n]$ with a linear order $\leq_{x}$ by setting $i \leq_{x} j$ if $x_{i} \leq x_{j}$. Then for each $i$, let $\sigma_{x}(i)$ denote the rank of $i$ in this order. Clearly $\sigma_{x}$ is a permutation of [ $n$ ], and every permutation can be realised in this way. Conversely, to each permutation $\sigma$ of $[n]$ we may associate the body $b_{\sigma}$ of all $x \in[0,1]^{[n]} \backslash \mathcal{B}$ such that $\sigma_{x}=\sigma$. Observe that (i) if $\sigma, \sigma^{\prime}$ are distinct elements of $S_{n}$, then $b_{\sigma}$ and $b_{\sigma^{\prime}}$ are disjoint subsets of $[0,1]^{[n]}$, and (ii) $\operatorname{vol}\left(b_{\sigma}\right)=1 / n!$ for all $\sigma \in S_{n}$. (It is also worth remarking that given a body $b \subseteq[0,1]^{V_{n}}$, we can define a $b$-random permutation by selecting $x \in b$ uniformly at random. This gives an interesting nonuniform model for random permutations.)

Given permutations $\sigma \in S_{N}$ and $\tau \in S_{n}$, we say $\sigma$ is a subpattern of $\tau$ if there is an orderpreserving injection $\phi:[N] \rightarrow[n]$ such that for every $i, j \in[N], \sigma(i)<\sigma(j)$ if and only if $\tau(\phi(i))<\tau(\phi(j))$. One important topic of study in permutation theory is that of pattern avoidance. Can one count or characterise the permutations in $S_{n}$ avoiding a given pattern $\sigma \in S_{N}$ ?

Definition 5.1. A permutation class $\mathcal{P}$ is a sequence $\left(\mathcal{P}_{n}\right)_{n \in \mathbb{N}}$ of subsets $\mathcal{P}_{n} \subseteq S_{n}$ which is closed under taking subpatterns. (i.e., if $\tau \in \mathcal{P}_{n}$ and $\sigma \in S_{N}$ is a subpattern of $\tau$ then $\sigma \in \mathcal{P}_{N}$.).

Definition 5.2. Given a permutation $\pi \in S_{N}$, let $S_{n}(\pi)$ denote the collection of all $\tau \in S_{n}$ that do not contain $\pi$ as a subpattern.

The Stanley-Wilf limit of the permutation class $\mathcal{S}=\left(S_{n}(\pi)\right)_{n \in \mathbb{N}}$ is

$$
L(\pi):=\lim _{n \rightarrow \infty}\left|S_{n}(\pi)\right|^{1 / n}
$$

(The existence of the limit $L(\pi)$ is highly nontrivial. Its existence was established by Marcus and Tardos in 2004 in [29].) We observe here that pattern-avoidance and Stanley-Wilf limits for permutation classes fit very nicely within the framework of [0, 1]-decorated ssee-s. To wit: let $\mathbf{V}$ be the ssee with $V_{n}=[n]$ and $\binom{V_{n}}{V_{N}}$ being the collection of all order-preserving injections $\phi: V_{N} \rightarrow V_{n}$. One can easily check that this constitutes a good ssee. Given a forbidden pattern $\pi \in S_{N}$, let $\mathcal{P}=\operatorname{Forb}\left(b_{\pi}\right)$. Clearly, we have $\left|S_{n}(\pi)\right|^{1 / n}=\left(n!\operatorname{vol}\left(\mathcal{P}_{n}\right)\right)^{1 / n}$. Thus providing a good estimate on $\operatorname{vol}\left(\mathcal{P}_{n}\right)$ via containers could potentially give a good estimate on $L(\pi)$. However, as pointed out in Section 5.4, the consequences of simple container theory obtained in this paper are not sufficiently precise to do so: In the language of graph theory, what we study corresponds to the "dense" case with strictly positive Turán density, whereas pattern avoidance belongs to the "sparse" case with zero Turán density.

The example of permutations suggests it would be interesting to obtain versions of Theorem 1.12 that work in a sparser setting, that is, with sharper estimates on the size of the container family $\mathcal{C}$ than are given by (iii). In this case, one will have to go back to the original theorems of Balogh-Morris-Samotij and Saxton-Thomason, rather than use the simple (but weaker) container theorem of Saxton-Thomason as a black box.

## 5.6 | Entropy maximisation in the decorated graph setting

Recall that the discrete entropy or Shannon entropy of a random variable $X$ taking values inside a discrete set $S$ is $\sum_{s \in s}-\mathbb{P}(X=s) \log (\mathbb{P}(X=s))$. The entropy we consider in this paper (see Definition 1.6) can be viewed as a continuous analogue of discrete entropy when $X$ is a point sampled uniformly at random from some body $b \in \mathbb{R}^{n}$.

In the $\{0,1\}$-decorated setting, the rough structure of discrete entropy maximisers for hereditary properties of graphs is well-understood, via the choice number $\chi_{c}$ (see [2, 8] for the set of possible "entropy densities" $\pi(\mathcal{P})$ and [4] for the possible structure of entropy maximisers). By contrast, it is less clear what the set of possible values of entropy densities or the possible rough structure of graphs maximising entropy should be in the [ $k$ ]-decorated setting for $k \geq 3$, let alone the set of entropy maximisers in the setting of [0, 1]-decorated graphs. We are only aware of one partial result in this area: Alekseev and Sorochan [3] who established a dichotomy on the growth rate of a symmetric hereditary property of [ $k$ ]-decorations of $E\left(K_{n}\right)$. Moreover, it is clear that the possible structures of entropy maximisers are much more varied than in the case $k=2$, see the discussion at the end of [18]. This leads to the following analytic problems.

Problem 5.3. Let $k \in \mathbb{N}$ with $k \geq 3$. Let $\mathcal{P}$ be a hereditary property of $[k]$-decorations of $E\left(K_{n}\right)$ and $\widehat{\mathcal{P}}$ be its completion under the cut norm. Determine the set of possible values for $m(\mathcal{P})=\sup _{W \in \widehat{\mathcal{P}}} \operatorname{Ent}(W)$, as well as the possible structures of entropy maximisers.

Problem 5.4. Let $\mathcal{P}$ be a hereditary property of [0, 1]-decorated graphs and $\widehat{\mathcal{P}}$ be its completion under the cut norm. Determine the set of possible values for $m(\mathcal{P})=\sup _{W \in \widehat{\mathcal{P}}} \operatorname{Ent}(W)$, as well as the possible structures of entropy maximisers.

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## DATA AVAILABILITY STATEMENT

The data that support the findings of this study are openly available in arXiv at https://arxiv. org/, reference number 2107.00550.

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