Reconciling SHACL and Ontologies: Semantics and Validation via Rewriting

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Abstract. OWL and SHACL are two prominent W3C standards for managing RDF graphs, the data model of the Web. They are used for different purposes and make different assumptions about the completeness of data: SHACL is used for expressing integrity constraints on complete data, while OWL allows inferring implicit facts from incomplete data; SHACL reasoners perform validation, while OWL reasoners do logical inference. Integrating these two tasks into one uniform approach is a relevant but challenging problem. The SHACL standard envisions graph validation in combination with OWL entailment, but it does not provide technical guidance on how to realize this. To address this problem, we propose a new intuitive semantics for validating SHACL constraints with OWL 2 QL ontologies based on a suitable notion of the chase. We propose an algorithm that rewrites a set of recursive SHACL constraints (with stratified negation) and an OWL 2 QL ontology into a stand-alone set of SHACL constraints that preserves validation for every input graph, which can in turn be evaluated using an off-the-shelf SHACL validator. We show that validation in this setting is EXP TIME-complete in combined complexity, but only PTIME-complete in data complexity, i.e., if the constraints and the ontology are fixed.

1 Introduction

SHACL and OWL are two prominent W3C standards for managing RDF data, the graph-based data model of the Web. They were specifically designed to target two different issues. OWL was standardized in parallel with RDF to address information incompleteness of RDF data by means of ontological axioms that complete the data with missing information. OWL and its profiles are based on Description Logics (DLs) and make the open-world assumption (OWA), which intuitively means that the data only presents a partial description of the domain of interest and missing facts may also be true. RDF and OWL were soon adopted by increasingly many applications, including enterprise systems and critical areas such as healthcare, and making decisions based on correct data became particularly crucial.

To check the correctness of RDF data, W3C proposed the so-called Shapes Constraint Language (or SHACL) [17], a machine-readable constraint language for describing and validating RDF graphs. Unlike OWL, SHACL operates under the closed-world assumption (CWA) and assumes completeness of data. SHACL specifies the notion of a shapes graph, which consists of a set of shape constraints paired with the so-called targets, which is a selection of nodes of the data graph that must be validated against the constraints. The precise semantics of SHACL in the presence of recursion was not described in the W3C standard, which led to recent works that propose semantics based on first-order logic and logic programming [7, 15, 2].

Combining SHACL and OWL into a setting that allows to perform RDF data validation while taking into account the implicit facts inferred using an OWL ontology is a relevant but challenging problem. Indeed, the W3C SHACL specification envisions graph validation in the presence of OWL entailment but does not provide guidance on how to realize this.

To our knowledge, this has only been addressed in [16], which considers positive SHACL constraints only.

To see the benefits of taking into account ontologies when performing validation, consider an example of a simplistic database of pet owners containing the facts hasWingedPet(linda, blu), Bird(blue), PetOwner(john), hasPet(john, ace) and consider the simple constraint

\[
\text{petOwnerShape} \leftarrow \text{PetOwner} \lor \exists \text{hasPet}
\]

with the target petOwnerShape(linda), which asks to verify whether linda is a pet owner. Clearly, one would expect the input data to validate linda as a pet owner given that she has a winged pet. However, this is not the case since the setting is missing the background knowledge that owning a winged pet implies owning a pet. The latter can be expressed through the ontological axiom hasWingedPet ⊆ hasPet, which would allow us to obtain the desired validation result.

Identification of a proper semantics in this setting requires integrating the OWA of OWL and the CWA of SHACL. There are several proposals by the DL and database communities to relax the OWA and combine it with CWA [11, 14, 13, 10]. Another challenge when defining a validation semantics is dealing with the non-monotonic behavior of SHACL constraints due to the presence of negation. Roughly speaking, adding facts to the input data graph may cause a previously valid setting to become invalid. Such non-monotonic behavior is known when combining ontologies and negation in the so-called conjunctive queries or database constraints (see e.g., [5, 14, 8]).

The contributions of this paper can be summarized as follows:

- We present a novel notion of SHACL validation in the presence of a DL-Lite\(_\text{\textcopyright}\) ontology, the logic underlying OWL 2 QL [12]. Specifically, we consider stratified SHACL constraints, which support a limited form of recursion (specifically, limiting the interaction between recursion and negation). Our notion of stratification is derived...
The impossibility of such a rewriting for SHACL with negation given in Theorem 1 of [16] does not apply to our semantics, nor to the minimal-model semantics adopted in [16], as acknowledged by the authors in personal communication.

A basic concept is a concept name \( A \in N_C \) or an expression of the form \( \exists r \) such that \( r \in N_r^+ \). A concept inclusion is an expression of the form \( C \sqsubseteq D \) or \( C \sqsubseteq \neg D \), where \( C, D \) are basic concepts. A role inclusion is an expression of the form \( r \sqsubseteq s \) or \( r \sqsubseteq \neg s \), where \( r, s \) are roles. A TBox \( T \) is a set of concept and role inclusions. We use \( K \) to denote a pair \((T, A)\) of a TBox and an ABox. The semantics is defined in terms of interpretations \( I \) in the usual way. The notions of satisfaction and of model of an inclusion, a TBox, and an ABox are standard, as is the definition of logical entailment. In particular, we write \( T \models \gamma \) if every model of the TBox \( T \) is also a model of \( \gamma \) (where the latter may be an inclusion, a TBox, or an ABox). We say that \( A \) is consistent with \( T \) (or \( K \) is consistent) if there is a model of \( A \) and \( T \).

### Basic SHACL notions.

We let \( N_S \) denote a countably infinite set of shape names, disjoint from \( N_C, N_R, \) and \( N_I \). A shape expression \( \varphi \) is an expression built according to the following grammar:

\[
\varphi ::= c | A | \neg A | s | \neg s | \varphi \land \varphi | \exists r.\varphi,
\]

where \( c \in N_I, A \in N_C, s \in N_S, \) and \( r \in N_r^+ \). In the following, we use \( \exists r \) as a shorthand for \( \exists r.\top \). A shape expression \( \varphi \), a shape assignment \( S \) and an interpretation \( I \) with \( N_I \subseteq \Delta^I \), we define the set \( \varphi^{T,S} \subseteq \Delta^I \) inductively as shown in Figure 1.

Given an interpretation \( I \) and a shapes graph \((C, G)\), we are interested in validating the targets in \( G \) given the information in \( I \) and the constraints in \( C \). Specifically, the goal of validation is to find a shape assignment \( S \) with \( G \subseteq S \) and such that \( S \) satisfies \( C \) and \( I \). For recursive SHACL, there are several validation semantics, i.e., different meanings of “satisfy”. For illustration, we recall the supported model semantics from [7]. We say \( S \) is a supported model of \( C \) and \( I \), if the following holds: \( s(a) \in S \) iff there exists \( s \leftarrow c \in C \) with \( a \in \varphi^{T,S} \). Validation under the supported model semantics consists of finding a supported model \( S \) for \( I \) and \( C \) such that \( G \subseteq S \). The supported model semantics suffers from unfounded justifications. Consider the singleton constraint set \( C = \{ s \leftarrow \exists r.\varphi \} \), the target set \( G = \{ s(a) \} \), an interpretation \( I \) with \( (a, a) \in r^I \), and a shape assignment \( S = (a) \). Then \( S \) is a supported model of \( I \) and \( C \) and it validates \( s(a) \). However, the validation of \( s(a) \) is clearly not well-founded. Another semantics is the stable model semantics introduced in [2], which allows to eliminate such unfounded inferences.
3 Stratified SHACL

In this paper, we are interested in shape assignments where each shape atom has a proper justification. We do not resort to the full stable model semantics, as we are interested in cases with polynomial time data complexity. We focus on stratified sets of constraints which support negation, but also admit a partition of constraints (stratification) such that a justified shape assignment can be constructed by processing each partition individually. This approach is based on the syntax and semantics of stratified logic programs, and thus we mainly only adapt the notions from [3].

Definition 3.1. We say a shape name $s$ occurs negatively in a shape constraint $s' \iff \phi$ if $\neg s$ occurs in $\phi$. We say a shape name $s$ is defined in a set $C$ of constraints if $s \iff \phi \in C$ for some $\phi$.

A set $C$ of constraints is stratified if it can be partitioned into sets $C_0, \ldots, C_k$ such that, for all $0 \leq i \leq k$, the following hold:

1. If $i < k$ and $s'$ occurs in $\phi$ for some $s' \iff \phi \in C_i$, then $s'$ is not defined in $C_{i+1} \cup \ldots \cup C_k$.
2. If $s'$ occurs negatively in $\phi$ for some $s \iff \phi \in C_i$, then $s'$ is not defined in $C_i \cup \ldots \cup C_k$.

A set of constraints is stratified if it admits a stratification.

We now define the notion of a consequence operator $T_{\mathcal{I}, \mathcal{C}}$ that, given a shape assignment $S$, adds new shape atoms to satisfy the constraints that are fired by the constraints in $C$ based on $S$ and $\mathcal{I}$.

Definition 3.2. For a set $C$ of constraints, an interpretation $\mathcal{I}$ with $\mathcal{N}_I \subseteq \Delta^I$, we define an immediate consequence operator $T_{\mathcal{I}, \mathcal{C}}$ that maps shape assignments to shape assignments as follows:

$$T_{\mathcal{I}, \mathcal{C}}(S) = S \cup \{s(a) \mid s \iff \phi \in C \text{ and } a \in (\phi)^{I,S}\}.$$ 

We further let:

- $T_{\mathcal{I}, \mathcal{C}}(S) = S$.
- $T_{\mathcal{I}, \mathcal{C}}(S) = T_{\mathcal{I}, \mathcal{C}}(\mathcal{I}, \mathcal{C})^n(S)$.
- $T_{\mathcal{I}, \mathcal{C}}(S) = \bigcup_{n=0}^{\infty} T_{\mathcal{I}, \mathcal{C}}(S)$.

The following two propositions are a direct consequence of the characterizations from [3] in the context of stratified logic programs.

Proposition 3.3. If $C$ is a constraint set that does not define any shape names that occur negatively in $C$, then the following hold:

1. $T_{\mathcal{I}, \mathcal{C}}$ is monotonic, i.e. if $S \subseteq S'$, then $T_{\mathcal{I}, \mathcal{C}}(S) \subseteq T_{\mathcal{I}, \mathcal{C}}(S')$;
2. $T_{\mathcal{I}, \mathcal{C}}$ is finitary, i.e. $T_{\mathcal{I}, \mathcal{C}}(\bigcup_{n=0}^{\infty} S_n) \subseteq \bigcup_{n=0}^{\infty} T_{\mathcal{I}, \mathcal{C}}(S_n)$ for all infinite sequences $S_0 \subseteq S_1 \subseteq \ldots$;
3. $T_{\mathcal{I}, \mathcal{C}}$ is growing, i.e. $T_{\mathcal{I}, \mathcal{C}}(S_2) \subseteq T_{\mathcal{I}, \mathcal{C}}(S_1)$ for all $S_1, S_2, S_3$ such that $S_1 \subseteq S_2 \subseteq S_3 \subseteq \mathcal{T}_{\mathcal{I}, \mathcal{C}}$.

Proposition 3.4. If $\mathcal{I}$ is an interpretation and $C_0, \ldots, C_k$ is a stratification of $C$, then all $T_{\mathcal{I}, \mathcal{C}_0}, \ldots, T_{\mathcal{I}, \mathcal{C}_k}$ are monotone, finitary, and growing. Thus, for any shape assignment $S$ and each $0 \leq j \leq k$, $T_{\mathcal{I}, \mathcal{C}_j} = T_{\mathcal{I}, \mathcal{C}_j}^n(S)$ is a fixpoint of $T_{\mathcal{I}, \mathcal{C}_j}$.

Based on the above, we can now define the computation of the desired shape assignment along a stratification $C_0, \ldots, C_k$ of $C$.

Definition 3.5. Assume $\mathcal{I}$ is an interpretation, $C$ is a stratified set of constraints, and let $C_0, \ldots, C_k$ be a stratification of $C$. Then let

- $M_0 = T_{\mathcal{I}, \mathcal{C}_0}^n(\emptyset)$
- $M_i = T_{\mathcal{I}, \mathcal{C}_i}^n(M_{i-1})$ for each $1 \leq i < k$.

We let $M_k$ be the perfect assignment for $C$ and $\mathcal{I}$, and let $PA(\mathcal{C}, \mathcal{I}) = M_k$. An interpretation $\mathcal{I}$ (resp., ABox $\mathcal{A}$) validates a shapes graph $(C, \mathcal{G})$ if $\mathcal{G} \subseteq PA(C, \mathcal{I})$ (resp., $\mathcal{G} \subseteq PA(\mathcal{C}, \mathcal{I}_A)$).

Assume an interpretation $\mathcal{I}$ and a shapes graph $(C, \mathcal{G})$, where $C$ is a stratified set of constraints. We first note that $PA(C, \mathcal{I})$ does not depend on a concrete stratification, i.e. any stratification yields the same $PA(C, \mathcal{I})$. Moreover, $PA(C, \mathcal{I})$ is a supported model of $C$ and $\mathcal{I}$, which follows from the corresponding result in [3] for stratified logic programs. Since the stable model semantics for SHACL in [2] was defined for finite structures only, we cannot precisely formalize here a connection between our semantics and the semantics in [2].

4 Validating SHACL with Ontologies

In this section we propose an intuitive validation semantics for SHACL in the presence of a DL-Lite$^\mathcal{K}$ ontology. More precisely, for a given a TBox $\mathcal{T}$, an ABox $\mathcal{A}$, and a shapes graph $(C, \mathcal{G})$, we need to define when $(\mathcal{T}, \mathcal{A})$ validates $(C, \mathcal{G})$. A natural first idea would be to follow the usual open-world semantics of DL-Lite and check for validation over all models of $\mathcal{A}$ and $\mathcal{T}$. While this works for positive constraints, it does not yield a natural result in the presence of negation, as illustrated in the following simple example.

Example 4.1. Consider an ABox $\mathcal{A}$, consisting of the facts hasPet(linda, blu), Bird(blue), PetOwner(linda), and an empty TBox $\mathcal{T}$. Let $(C, \mathcal{G})$ be a shapes graph, where $C$ only contains the constraint

$$s \iff \exists s\text{hasPet} \land \neg \exists s\text{hasPet} .$$

and the target to be checked for validation is $\mathcal{G} = \{s(linda)\}$.

The perfect assignment $PA(C, \mathcal{I}_A)$ naturally includes the target $s(linda)$ since linda has a pet, and does not have a pet that is a dog. Note that since we have an empty TBox, we are in the usual setting of validation here. Thus, $\mathcal{A}$ validates $(C, \mathcal{G})$ and clearly, one would expect $(\mathcal{T}, \mathcal{A})$ to validate $(C, \mathcal{G})$. However, if we consider all possible models of $(\mathcal{T}, \mathcal{A})$, we have non-validation since there may be a model $\mathcal{I}$ of $\mathcal{T}$ and $\mathcal{I}$ that includes a hasPet-fact for linda to some pet $b$ that is a Dog. Indeed, the interpretation $\mathcal{I}$ does not validate the shapes graph since the target $s(linda)$ is not included in $PA(C, \mathcal{I})$.

The problem in the above example is the non-monotonicity of the constraints. Roughly speaking, adding facts to the data may cause a previously valid setting to become invalid. We want an intuitive semantics that coincides with the usual validation in case the TBox is empty. As done in related settings (see e.g., [5, 14, 8]) we rely on the chase procedure [1] known from Knowledge Representation and Database Theory. Roughly speaking, a chase procedure takes as input an ABox and a TBox and iteratively applies the axioms of the TBox to the data by adding atoms over possibly fresh individuals until all the axioms in the TBox are satisfied. The result of the chase is a so-called canonical or universal model, and since it can be homomorphically embedded into every other model of the ABox and
TBox, the canonical model can be used as a representative of all the models. For DL-LiteR ontologies, such chase procedures may not terminate and result in infinite models. There are several chase variants producing different canonical models [8]. While for positive constraints these differences do not matter, constraints with negation can distinguish between them, resulting in different validation answers as we illustrate in Example 4.6.

### 4.1 Semantics over the Austere Canonical Model

The semantics we propose is based on a special chase procedure that constructs an austere canonical model, in the sense that each chase step introduces as few fresh successors as possible, without merging any successors unless the TBox is forcing us to do so.

We first define an auxiliary notion of a good successor configuration, which, given a node e and the set of basic concepts that are assumed to hold at e, determines the minimal set of fresh successors of e and the roles connecting e to each of them.

**Definition 4.2.** For a set of roles R, let $cl_T(R)$ be the set of all roles $r'$ such that $T \models r \subseteq r'$ and $r \in R$. Given a DL-LiteR TBox $T$ and a set of basic concepts $U$, a good successor configuration $succ_T(U)$ for $U$ is a possibly empty set of roles such that:

1. There do not exist $r, r' \in succ_T(U)$ such that $r \in cl_T(\{r'\})$;
2. If $C \in U$, $T \models C \subseteq \exists r$ and $\exists r' \notin U$, then there exists $r' \in succ_T(U)$ such that $r \in cl_T(\{r'\})$;
3. If $r \in succ_T(U)$, there exists $C \in U$ such that $T \models C \subseteq \exists r'$;
4. $succ_T(U) \cap cl_T(\{r \mid \exists r \in U\}) = \{\}$.

**Lemma 4.3.** For each set of basic concepts $U$, there exists a unique good successor configuration $succ_T(U)$ for $U$.

The following example illustrates the notion of a good successor configuration.

**Example 4.4.** Consider the TBox $T$ containing three axioms:

\[
P etOwner \sqsubseteq \exists \text{hasPet} \quad \text{hasWingedPet} \sqsubseteq \text{hasPet}
\]

Then, $succ_T(\{\text{PetOwner}\})$ is the set $\{\text{hasWingedPet}\}$, which trivially satisfies items (1) to (4). Indeed $succ_T(\{\text{PetOwner}\})$ cannot contain $\text{hasPet}$ since by item (1), $\text{hasPet} \in cl_T(\{\text{hasWingedPet}\})$.

All other sets of basic concepts will have empty sets as successor configurations. In particular, the set $U = \{\text{PetOwner}, \exists \text{hasWingedPet}\}$ cannot have any of the roles of the TBox as a successor role due to item (4), since $cl_T(\{\text{hasWingedPet}\}) = \{\text{hasWingedPet}, \text{hasPet}\}$.

The good successor configuration prescribes how to satisfy the TBox axioms locally. We use it to build a canonical model layer by layer that we call austere.

**Definition 4.5.** Let $K = (T, A)$ be a pair of a DL-LiteR TBox $T$ and ABox $A$. Let $N_I(A)$ be the set of individuals occurring in $A$. Let $N_K$ be the set of finite words of the form $ar_1r_2 \ldots r_n$, with $n \geq 1$, $a \in N_I(A)$ and $r_1, \ldots, r_n$ are roles, such that the following hold:

1. Let $U_a = \{A \mid A(a) \in A \} \cup \{\exists r \mid r(a, a') \in A\}$, then $r_1 \in succ_T(U_a)$;
2. For every $1 \leq i < n$, $r_{i+1} \in succ_T(U_{r_i})$, where $U_{r_i} = \{\exists r_i^+\}$.

We denote by $tail(w)$ the last role appearing in a word $w \in N_K$.

The **austere canonical model** $can(K)$ of $K$ is the interpretation $I$ with domain $\Delta_I := N_I(A) \cup N_K$ such that for all $a \in N_I(A)$, concept names $A$ and roles $r$, the following hold:

1. $a_I := a$;
2. $A_I := \{a \mid K \models A(a)\} \cup \{w \in N_K \mid T \models \exists tail(w) \subseteq A\}$;
3. $r_I := \{(a, a') \mid K \models r(a, a')\} \cup \{(w_1, w_2) \mid w_1 \in N_K, w_2 = w_1r', r \in cl_T(\{r'\})\}.$

For a consistent $K = (T, A)$, the austere canonical model $can(K)$ exists and is unique. It is standard to show that $can(K)$ is a canonical model of $K$, which is a model of $K$ that can be homomorphically embedded into every other model of $K$. The number of successor nodes in $can(K)$ is determined by the good successor configuration.

Our notion of austere canonical model is closely related to the core chase [8]. It will typically create fewer fresh successors than the **oblivious chase**, which, roughly speaking, applies the axioms of the TBox without first checking whether the axiom is already satisfied. It may also create fewer successors than the **restricted chase**, which may be sensitive to the order of rule applications.

**Example 4.6.** Consider $K = (T, A)$ where the ABox $A = \{\text{PetOwner}(\text{linda}), \text{hasWingedPet}(\text{linda}, \text{blu}), \text{Bird}(\text{blu})\}$ and the TBox $T$ as in Example 4.4. The good successor configuration $succ_T(\{\text{PetOwner}\})$ is the empty set. More precisely, $U_{\text{PetOwner}} = \{\text{PetOwner}, \exists \text{hasWingedPet}\}$ and as shown in Example 4.4, it cannot have any of the roles of the TBox as a successor role. Thus, $\text{N}_K$ in the canonical model is empty, and no fresh objects are introduced. The austere canonical model $can(K)$ (right in Figure 2) will only add a $\text{hasWP}$-role from $\text{linda}$ to $\text{blu}$. In contrast, the canonical model obtained from the oblivious chase (left in Figure 2) will introduce two fresh objects to satisfy the two existential axioms.

![Figure 2: Result of oblivious chase (left) and austere canonical model (right). We use hasWP and hasP instead of hasWingedPet and hasPet.](image)

Consider now the shapes graph $(C, G)$ with $C = \{s \leftarrow \exists \text{hasPet}.\neg \text{Bird}\}$ and $G = \{s(\text{linda})\}$. The shapes graph asks to validate whether $\text{linda}$ has a pet that is not a bird. Clearly, the austere canonical model provides the expected answer, as it does not validate $(C, G)$. In contrast, the canonical model on the left-hand-side of the figure—and thus the semantics of [14] adopted for SHACL in [16]—results in the unintended validation of $(C, G)$.

The semantics of validation with DL-LiteR ontologies is given in terms of validation over the austere canonical model.

**Definition 4.7 (Validation with DL-LiteR).** Assume $K = (T, A)$ and $(C, G)$ is a shapes graph, where $C$ is a stratified set of constraints. We say $K$ validates $(C, G)$ if $can(K)$ validates $(C, G)$.

We illustrate the notion of validation with an example.

**Example 4.8.** Consider again Example 4.6 and the shapes graph $(C, G)$ with $C = \{s \leftarrow \exists \text{hasPet}.\neg \text{Bird}\}$ and $G = \{s(\text{linda})\}$. 

![Figure 2: Result of oblivious chase (left) and austere canonical model (right). We use hasWP and hasP instead of hasWingedPet and hasPet.](image)
5 Rewriting

In the rest of the paper, we will only consider ABoxes that are consistent with the given TBox. Furthermore, note that disjointness axioms will not play any role in the rewriting.

Given a TBox $\mathcal{T}$ and a set of stratified constraints $C$, we want to compile $\mathcal{T}$ and $C$ into a new set $C_{\mathcal{T}}$ of stratified constraints so that for every ABox $\mathcal{A}$ consistent with $\mathcal{T}$, and every target $\mathcal{G}$, we have $(\mathcal{T}, \mathcal{A}) \text{ validates } (\mathcal{C}, \mathcal{G}) \iff \mathcal{A} \text{ validates } (\mathcal{C'}, \mathcal{G})$.

This will be achieved by means of an inference procedure that uses a collection of inference rules to capture the "possibility" of shape names in the anonymous part of the canonical model. To ease presentation, we assume $C$ contains only normalized constraints of the following forms:

\[(NC1) \ s \leftarrow D \quad (NC2) \ s \leftarrow s_1 \land s_2 \quad (NC3) \ s \leftarrow \exists r.s' \quad (NC4) \ s \leftarrow \neg s'\]

where $D$ is a disjunction or a (DL-Lite$\mathbb{R}$) basic concept. Every set of constraints can be easily normalized without affecting validity (see e.g., [2]). This can be done in polynomial time by introducing fresh shape names. Note that we may say negative constraints to refer to constraints of the form (NC4) and positive constraints to refer to constraints of the form (NC1) to (NC3).

In the first stage, we define the rewriting procedure for constraints without negation. We later show how that procedure can be upgraded to handle stratified negation.

5.1 Rewriting Positive Constraints

Our rewriting algorithm uses pairs of the form $(W, H)$, where $W$ is a set of basic shape expressions and $H$ is a set of shape names.

Definition 5.1. If $A$ is a concept name, $r$ is a role, $s$ is a shape name, and $c$ an individual, then $c, A, \exists r. T, \exists r.s, \neg c, \neg A, \neg \exists r.T$, and $\neg \exists r.s$ are called basic shape expressions; the former four are positive, the latter four are negative. For a set $W$ of basic shape expressions, we let

\[cr(W) = \{\exists r \mid \exists r.T \in W \text{ or } \exists r.s \in W\} \cup \{A \mid A \in W\}.\]

The core of our technique is an inference procedure that derives a set $I$ of pairs $(W, H)$. Intuitively, $(W, H) \in I$ tells us that an object that satisfies all expressions in $W$ validates all shape names in $H$.

Example 5.2. Let $\mathcal{T} = \{q \subseteq \exists p, \exists p^{-} \subseteq A\}$ and $C = \{s \leftarrow \exists p.s_A, s_A \leftarrow A, s' \leftarrow q\}$.

We want our rewritten $C_{\mathcal{T}}$ to expand the constraints $C$ above with

\[s \leftarrow q \quad s_A \leftarrow \exists p^{-}\]

allowing us to validate the target $s(a)$ over the ABox $\mathcal{A} = \{q(a, b)\}$, or the target $s_A(b)$ over the ABox $\mathcal{A} = \{p(a, b)\}$, without having to build $can(\mathcal{T}, \mathcal{A})$. This will be achieved by deriving a set of pairs that includes $\{(q\exists q), (s, s')\}$, which keeps the constraint $s' \leftarrow q$ while also adding the new constraint $s \leftarrow q$, as well as the pairs $\{\exists p^{-}, s_A\}$ and $\{A, s_A\}$, witness the constraints for $s_A$. In the following definition, we assume a given set $I$ of pairs $(W, H)$ as above. The role of this set will become clear later when we lift the algorithm to constraints with stratified negation.

Definition 5.3. Let $\mathcal{T}$ be a TBox and $C$ a set of normalized constraints. Let $\Sigma_{\mathcal{C}, \mathcal{T}}$ be the set of all basic shape expressions that can be formed from the concept, role, and shape names occurring in $\mathcal{T}$ and $C$. We let $sat_{\mathcal{C}, \mathcal{T}}(I)$ be the smallest set of pairs, containing $I$, closed under the following rules:

1. If $s \in \Sigma_{\mathcal{C}, \mathcal{T}}$ is a positive basic shape expression, then $\{(S), \{\}\}$ belongs to $sat_{\mathcal{C}, \mathcal{T}}(I)$.
2. If $s \leftarrow s' \in C$ for some basic shape expression $S$, then $\{(S'), \{s\}\}$ belongs to $sat_{\mathcal{C}, \mathcal{T}}(I)$ for each $S' \in \Sigma_{\mathcal{C}, \mathcal{T}}$ such that:
   - $S' = S = c$ for some individual $c$;
   - $\mathcal{T} \models s' \subseteq S$ and $S'$ are basic DL-Lite$\mathbb{R}$ concepts; or
   - $\mathcal{T} \models r' \subseteq r$ and $S, S'$ are of the forms $\exists r.s'$ and $\exists r'.s'$ respectively.
3. If $\{(W, H), (W', H')\} \subseteq sat_{\mathcal{C}, \mathcal{T}}(I)$ such that there is no basic concept or individual $D$ with $\{D, \neg D\} \subseteq W \cup W' \cup cr(W) \cup cr(W')$, and no two distinct individuals $c, c' \subseteq N_1$ with $\{c, c'\} \subseteq W \cup W'$, then $(W \cup W', H \cup H')$ belongs to $sat_{\mathcal{C}, \mathcal{T}}(I)$.
4. If $s \leftarrow s_1 \land s_2 \in C$, $(W, H) \in sat_{\mathcal{C}, \mathcal{T}}(I)$ and $s_1, s_2 \subseteq H$, then $(W, H \cup \{s\}) \in sat_{\mathcal{C}, \mathcal{T}}(I)$.
5. If $s \leftarrow \exists r.s' \in C$, $(\{W, H\}, \{W', H'\}) \subseteq sat_{\mathcal{C}, \mathcal{T}}(I)$, $s' \in H'$ and there are $r' \in N_1 R$ and $R \subseteq N_1 R$ such that:
   - $r' \in succ_{\mathcal{T}}(cr(W))$, $r = cl_{\mathcal{T}}(\{r'\})$, and $r \in R$,
   - $cr(W') = \{\exists r \mid r \subseteq R\}$,
   - $\{s'' \mid \exists r'.s'' \in W'\} \cup \{\neg s'' \mid \neg \exists r'.s'' \in W', r'' \subseteq R\} \subseteq H$,
   - then $(W \cup \{\exists p \mid r \in cl_{\mathcal{T}}(\{p\})\}, H \cup \{s\}) \in sat_{\mathcal{C}, \mathcal{T}}(I)$.

We let $C_{\mathcal{T}, I}$ be the set that contains, for each $(W, H) \in I$ and each $s \in H$, the constraint

\[s \leftarrow \bigwedge_{s \in W} S.\]

We let $C_{\mathcal{T}} = C_{\mathcal{T}, I}$ with $I = sat_{\mathcal{C}, \mathcal{T}}(\emptyset)$.

The following example illustrates the application of each of the rules in Definition 5.3.

Example 5.4. Consider $\mathcal{T} = \{A \subseteq \exists p, \exists p^{-} \subseteq \exists q, \exists q^{-} \subseteq B, p \subseteq r, q \subseteq r\}$ and $C = \{s \leftarrow \exists r.s, s \leftarrow s_B \land s', s_B \leftarrow B, s' \leftarrow \exists r^{-}.s', s \leftarrow A\}$.

Given the ABox $\mathcal{A} = \{A(a)\}$, the austere canonical model $can(\mathcal{T}, \mathcal{A})$ can be found in Figure 3.

\[a: A \quad \leftarrow \quad p, r \quad \leftarrow \quad b \quad \quad q, r \quad \leftarrow \quad c, B\]

Figure 3: Austere canonical model $can(\mathcal{T}, \mathcal{A})$ from Example 5.4

Let us consider the target $\mathcal{G} = \{s(a)\}$. It is now an easy check that $can(\mathcal{T}, \mathcal{A})$ validates $(C, \mathcal{G})$: just compute the perfect assignment.

To see that we can reach the same result with our rewriting, first apply Rule 2 two times to find $\{\{\exists q^{-}\}, \{s_B\}\}, \{\exists r^{-}.s', \{s'\}\} \subseteq$
I, where we use $I$ as a shorthand for $satc_{T}(\emptyset)$. By applying first Rule 3, then Rule 4 on these two pairs, we find that \((\exists p^r, \exists r^s, s), (s_{B}, s', s)\) \in I too. Call this pair \((W', H')\). Note that this pair is simulating the environment of $C$, which has as incoming $g$ edge, and an incoming $r$ edge from $b$, for which we assume that $s'(b)$ is validated. In such a context, we can correctly infer that $s_B(c), s'(c)$, and $s(c)$ are validated. Let’s now look at the environment $W$ of $b$, which contains $\exists p^r$ and $\exists r^s$, and, according to the good successor configuration, such a node has a $q$-successor.

Let \((W, H) = ((\exists p^r, \exists r^s), s')\) be constructed from applying Rule 2 on $s' \leftarrow \exists r - s'$. Intuitively, this reflects that $s(a)$ and $s'(b)$ are in the perfect assignment), followed by Rule 1 and 3 to add $\exists p^r$ to $W$. We can now apply Rule 5 where we consider the constraint $s' \leftarrow \exists q - s':$ let $r' = q, R = \{q, r\}$. Since \((W', H')\) is simulating the environment of $C$ and $\exists r^s, s' \in W'$, we need to check that $s' \in H$, or in words, that $s'(b)$ is validated. It is now easy to see that all conditions of Rule 5 are met: \((\exists p^r, \exists r^s, s' - \exists q^s), (s', s)\) \in I. Note that the function of $\exists q^s$ is to make sure that we can never add $\exists q$ to the left-hand side of the pair; adding this would break our assumption that a $q$-successor is introduced by the good successor configuration.

Next, we apply Rule 5 on the pair \((W, H) = ((A), s')\) and the pair \((W', H') = ((\exists p^r, \exists r^s), \{s', s\})\). We find that \((A, \exists p^r), \{s', s\}\) \in I. This means that the following constraint is in $C_{T}$:

$$s \leftarrow A \land \neg \exists p,$$

which shows that indeed $A$ validates \((C_T, G)\).

We generalize this in the following theorem: the rewriting gives a sound and complete set of constraints.

**Theorem 5.5.** Let $T$ be a TBox and $C$ a set of positive constraints. Then for every target $G$ and every ABox $A$ that is consistent with $T$, we have that $(T, A)$ validates $(C, G)$ iff $A$ validates $(C_T, G)$.

**Proof.** (Sketch) To show soundness, assume an arbitrary but fixed sequence of rule applications leading to $satc_{T}(\emptyset)$, let $I_0$ be the set of pairs after $i$ steps, and let $C_{T, i}$ be the constraints obtained from $I_i$. It suffices to show by induction on $i$, that for every arbitrary target $s(a)$, if $can(T, A)$ validates $(C_{T, i}, \{s(a)\})$, then $can(T, A)$ validates $(C, \{s(a)\})$. For $i = 1$, only Rule 1 or 2 can be applied. In the case of Rule 1 the claim is trivial as pairs $(W, \emptyset)$ do not generate constraints. In the case it is Rule 2, it is also straightforward since the added $(W, H)$ will result in a constraint of the form $s \leftarrow S'$ and if $S'$ is a basic concept then there is some basic concept $S$ such that $s \leftarrow S \in C$, and from $T \models S' \subseteq S$ it is the case that $a \in S'_{can(T, A)}$ implies $a \in S_{can(T, A)}$. Hence, if $s(a) \in PA(C_{T, i}, can(T, A))$, then $s(a) \in PA(C, can(T, A))$. The argument for $S'$ of the forms $\exists r.s$ and $\exists r'.s'$ is similar. For the induction step, by inspecting the Rules 2 to 5 one can easily show that each pair $(W, H)$ is correctly obtained at rule application $i + 1$ and it will generate sound constraints.

For completeness, suppose $can(T, A)$ validates $(C, \{s(a)\})$. Then there exists a justification tree indicating which domain elements of $can(T, A)$ validate which shape names that are used in the fixed point computations. In this tree, nodes are of the form $(x, S)$, where $x \in N_{I}(A) \cup N_{K}$ and $S \in N_{S} \cup \{D \mid D$ a basic concept$\}$. The root is $(a, s)$. Each node has none, one or two children and there are no infinite branches. Moreover, the following conditions hold:

- $(x, S)$ has no children if $S$ is a basic concept. If $(x, S)$ appears as a leaf and $S = A$ for some concept $A$, then $x \in A_{can(T, A)}$; and if $S = \exists r$ for some role $r$, there exists $y \in D_{can(T, A)}$ such that $(x, y) \in r_{can(T, A)}$.
- If $(x, S)$ has only one child, if it is of the form $(x, D)$ or $(x', s')$. If the form is $(x, D)$, then $s \leftarrow D \in C$. If the form is $(x', s')$, then $x' = x \cdot r'$ such that $r \in cl_{T}(\{r'\})$ or $x = x \cdot r'$ such that $r \in cl_{T}(\{r'\})$ and $s \leftarrow \exists r.s' \in C$. If $(x, S)$ has two children, they are of the form $(x, s')$ and $(x, s'')$, and $s \leftarrow s' \land s'' \in C$.

Define the weight of a tree as the number of nodes $(w, S)$ with $w \in N_{C}$. We call a leaf strictly justified if there exists a pair $(W, H)$ such that $cr(W) = \{\exists r \mid r \in cl_{T}(\{tail(x)\})\}$ and $\exists r.s \in w$, the child node, as soon as we have it, into a justification for the parent. If on a branch of the tree there are no nodes such that $(x, S)$ is the parent of $(x', S')$ and $|x| > |x'|$, we call this branch a loose end. The weight can be lowered easily by cutting off leaves on loose ends.

However, a $U$-turn can also happen in the anonymous part of our justification: from $(x, S)$ there exists a path to $(x', S')$ such that all nodes appearing on this path are of the form $(x', S'')$, where $x' = x \cdot r$ for some $r$. If we are cutting off leaves on loose ends, we might find a node with two children, one of them a leaf. In this case, we cut off the leaf and mark the edge to the other child. The other option is that we find a connected set of nodes $M$ such that if $(x', S') \in M$, then $x' = x \cdot r$ and if some $(x'', S') \not\in M$ is the child or parent of some $(x', S') \in M$, then $x'' = x$. We can take out complete sets $M$ at once, in such a way that the parent of this set will become the new parent of $M$’s children, but with a marked edge. Note that rule 5 deals explicitly with this case.

As the weight of the final tree is 0, we can use Rules 1, 2 and 3 to create pairs $(W, H)$ for each leave node left. Normal edges correspond to already existing shapes that we can use here too. Marked edges tell us that there exists a rewriting from the child to the parent: we know we can rewrite the incoming pair $(W, H)$ in the needed form. This tells us how $A$ validates $(C_T, \{s(a)\})$.

**5.4 Constraints with Stratified Negation**

We now extend the above rewriting to additionally handle constraint sets $C$ that involve stratified negation. Intuitively, this is done by running the saturation procedure at each stratum of $C$, starting with the lowermost. For this, we need to make sure that the outcome of the saturation at a non-topmost stratum is completed with negative information and is passed to the next stratum. To this end, we w.l.o.g. assume that all constraints from $C$ with the same shape name on the left-hand-side occur together in the same stratum.

We will operate now on pairs $(W, H)$, which are similar to the ones in the previous section, except that $H$ might additionally contain expressions of the form $\neg s$ for a shape name $s$. For a set $I$ of such pairs, we say $(W, H)$ is maximal in $I$, if $(W, H) \in I$ and there is no $H' \supset H$ with $(W', H') \in I$ such that $W'$ contains exactly the same basic shape expressions as $W$. Then the notion of completion...
is defined as follows:

**Definition 5.6.** The completion \( \text{comp}_{\mathcal{C}, \mathcal{T}}(I) \) of a set \( I \) of implication pairs w.r.t. a TBox \( \mathcal{T} \) and a constraint set \( \mathcal{C} \) is a set defined as follows:

\[
\text{comp}_{\mathcal{C}, \mathcal{T}}(I) = \{(W \cup \overline{W}, H \cup \overline{H}) \mid (W, H) \text{ is maximal in } I\}
\]

where

\[
\overline{H} := \{\neg s \mid s \text{ occurs in } \mathcal{C}, s \not\in \mathcal{H}\}
\]

\[
\overline{W} := \{\neg D \mid \text{there is no } D' \in \mathcal{D}(W) \text{ s.t. } \mathcal{T} \models D' \subseteq D\} \cup \{\neg \exists r.s' \mid s' \in \mathcal{C}, s \not\in \mathcal{H}\} \cup \{\neg c \mid s \not\in \mathcal{C}, s \not\in \mathcal{H}\}.
\]

We need to augment the inference rules of the previous section with an additional rule to handle constraints of the form \( s \leftarrow \neg \exists'. \)

**Definition 5.7.** Let \( \mathcal{T} \) be a TBox, \( \mathcal{C} \) a set of normalized constraints. We let \( \text{sat}_{\mathcal{C}, \mathcal{T}}(I) \) be the smallest set of pairs containing \( I \) closed under Rules 1–5 in Def. 5.3, and additionally under the following Rule 6:

6. If \( s \leftarrow \neg s' \in \mathcal{C} \) and \( (W, H) \in \text{sat}_{\mathcal{C}, \mathcal{T}}(I) \) such that \( \neg s' \in H \), then \( (W, H \cup \{s\}) \) belongs to \( \text{sat}_{\mathcal{C}, \mathcal{T}}(I) \).

Now we can define the inference procedure that processes strata from the lowest to the topmost, performing saturation using the updated set of rules at every stratum, interleaved with a computation of the completion in between.

**Definition 5.8.** For a TBox \( \mathcal{T} \) and a constraint set \( \mathcal{C} \) with stratification \( C_0, \ldots, C_n \), we let \( I_0 = \text{sat}_{C_0, T}(\emptyset) \) and for \( 0 < i \leq n \)

\[
I_i = \text{sat}_{C_i, \mathcal{T}}(\text{comp}_{C_{i-1}, \mathcal{T}}(I_{i-1})).
\]

We let \( \mathcal{C}_T = \mathcal{C}_{T,n} \), where \( \mathcal{C}_{T,n} \) is defined as in Definition 5.3.

**Example 5.9.** Consider \( \mathcal{T} = \{A \sqsubseteq \exists p\} \) and the set \( \mathcal{C} \) of constraints:

\[
s \leftarrow \exists p.s' \quad s' \leftarrow \neg s_B \quad s_B \leftarrow B.
\]

Let \( C_0, C_1 \) be a stratification of \( \mathcal{C} \) such that \( C_0 = \{s_B \leftarrow B\} \) and \( C_1 = \{s \leftarrow \exists p.s', s' \leftarrow \neg s_B\} \). The rewriting, among others, will produce the pair \( (\exists p.s'), \{\} \) with Rule 1 which is maximal in \( I_0 \), and hence, the pair \((\exists p.s', \neg A, \neg B, \neg \exists p, \neg s_B)\) is included in \( \text{comp}_{C_0, \mathcal{T}}(I_0) \). By Rule 6, the pair \((W', H') = (\exists p.s', \neg A, \neg B, \neg \exists p, \neg s_B, s')\) is added in \( I_1 \). By applying Rule 5 to the pairs \((W', H')\) and \( (W, H) = (\{A\}, \{\}) \), the pair \((\{A, \neg \exists p\}, \{s\})\) is added to \( I_1 \). This is translated into the constraint \( s \leftarrow A \neg \exists p \in \mathcal{C}_T \). Consider now the ABox \( A_1 = \{A(a)\} \). Clearly, \( \{A_1, \mathcal{T}\} \) validates \( (C, \{s(a)\}) \) and \( \{A_1, \mathcal{C}_T, \{s(a)\}\} \).

The rewriting gives other pairs \((W, H) \in I_1\) as well, including \((\{B\}, \{s_B\}), (\exists p.s'), \{s\}) \) and constraints in which a subset of \( \{A, B, \exists p, \exists p, \neg s_B, s'\} \) is added to \( W \). In addition, by taking the completion of \( \text{sat}_{C_0, \mathcal{T}}(\emptyset) \), we find pairs with \( \neg s_B \in H \) if \( B \notin W \). Now consider the ABox \( A_2 = \{A(a), p(a, b), B(b)\} \). We cannot use any pair \((W, H) \) introduced by Rule 5, as \( W \) then contains \( \neg \exists p \), so we have that \( A_2 \) does not validate \( (\mathcal{C}_T, \{s(a)\}) \), as desired.

Now consider \( A_3 = \{A(a), p(a, b)\} \). \( A_3 \) validates \( \{C_T, \{s'(b)\}\} \) witnessed by \( (\exists p.s', \neg A, \neg B, \neg \exists p, \neg s_B, s') \in I_1 \), and it validates \( \{C_T, \{s(a)\}\} \), witnessed by \( (\exists p.s'), \{s\} \in I_1 \).

**Theorem 5.10.** Let \( \mathcal{T} \) be a TBox and \( \mathcal{C} \) a set of stratified constraints. Then, for every target \( \mathcal{G} \) and every ABox \( \mathcal{A} \) that is consistent with \( \mathcal{T} \), we have that \((\mathcal{T}, \mathcal{A}) \) validates \((\mathcal{C}, \mathcal{G}) \) iff \( \mathcal{A} \) validates \((\mathcal{C}_T, \mathcal{G}) \).
To prove \textsc{Exptime}-hardness in combined complexity we reduce the word problem of polynomially space-bounded Alternating Turing Machines (ATMs) to validating shapes under a DL-Lite\(_{r}\) ontology.

An ATM is defined as a tuple of the form

\[ M = (\Sigma, Q_\exists, Q_\forall, q_0, q_{acc}, q_{rej}, \delta) \]

where \( \Sigma \) is an alphabet, \( Q_\exists \) is a set of existential states, \( Q_\forall \) is a set of universal states, disjoint from \( Q_\exists \), \( q_0 \in Q_\exists \) is an initial state, \( q_{acc} \in Q_\exists \cup Q_\forall \) is an accepting state and \( q_{rej} \in Q_\exists \cup Q_\forall \) is a rejecting state, and \( \delta \) is a transition relation of the form

\[ \delta \subseteq Q \times (\Sigma \cup \{B\}) \times Q \times (\Sigma \cup \{B\}) \times \{-1, 0, +1\} \]

with \( Q = Q_\exists \cup Q_\forall \). Here \( B \) is the the blank symbol. We let \( \delta(q,a) = \{(q', b, D) \mid (q, a, q', b, D) \in \delta\} \). W.l.o.g. we assume that in \( M \), universal and existential states are strictly alternating: if \( (q, a, q', b, m) \in \delta \) and \( q \in Q_\exists \) (resp., \( q \in Q_\forall \)), then \( q' \in Q_\exists \) (resp., \( q' \in Q_\forall \)). We further assume that \( |\delta(q,a)| = 2 \) for all combinations of states \( q \in Q \) and symbols \( a \in \Sigma \). If \( \delta(q,a) = \{(q_1, a_1, D_1), (q_2, a_2, D_2)\} \), we let \( \delta_i(q,a) = (q_1, a_1, D_1) \) and \( \delta_i(q,a) = (q_2, a_2, D_2) \).

A run of an ATM \( M \) on an input word \( w \) is defined as usual. We assume a word \( w = d_1 \cdot \ldots \cdot d_n \in \Sigma^* \) with \( n > 0 \) together with an ATM \( M \) that only uses the tape cells where the input word was written, i.e., it only uses the first \( n \) cells. Checking if such \( M \) accepts \( w \) is an \textsc{Exptime}-hard problem.

We show how to construct \( K = (T, A) \) and \( (C, \{s(a)\}) \) such that \( M \) accepts \( w \) iff \( K \) validates \( (C, \{s(a)\}) \). The reduction takes polynomial time in the size of \( M \) and \( w \). It uses the following symbols:

- a concept name \textit{Init} and a shape name \textit{sacc};
- role names \textit{succ}, \textit{succ}_c, \textit{succ}_r;
- shape names \( s_q \) for all states \( q \in Q_\exists \cup q \in Q_\forall \);
- shape names \( h_i \) for all \( 1 \leq i \leq n \);
- shape names \( c_{i,j} \) for all \( i \in \Sigma \cup \{B\} \) and \( 1 \leq i \leq n \).

We let \( K = (T, A) \) be such that \( A = \{\text{Init}(a)\} \) and the TBox \( T \) consists of the following inclusions:

\[
\begin{align*}
\text{Init} & \sqsubseteq \exists \text{succ}_c \quad \text{succ}_c \sqsubseteq \text{succ} \quad \exists \text{succ}^- \sqsubseteq \exists \text{succ}_c \\
\text{Init} & \sqsubseteq \exists \text{succ}_r \quad \text{succ}_r \sqsubseteq \text{succ} \quad \exists \text{succ}^- \sqsubseteq \exists \text{succ}_r.
\end{align*}
\]

The interpretation can(\( K \)) will provide us an infinite binary tree. In there, the root is representing the starting cells where the input word was written and each child of a node represents a next step in the run of the ATM \( M \).

To mimic the start configuration, we define the following shapes:

\[
\begin{align*}
h_1 & \leftarrow \text{Init} \quad s_{q_0} \leftarrow \text{Init} \quad c_{i,i}^{(1)} \leftarrow \text{Init} \quad \text{for all } 1 \leq i \leq n
\end{align*}
\]

Intuitively, this is setting the starting state to \( q_0 (s_{q_0}) \), putting the head in the starting position (\( h_1 \)), and starting the shape written on the tape cell (\( c_{i,i}^{(1)} \)).

The next step is to encode the transition relation of the \( M \). For each \( 1 \leq i \leq n \), each \( (q, a) \in Q \times (\Sigma \cup \{B\}) \), and \( \gamma \in \{\ell, r\} \) we add the following shapes, where \( (q', b, D) = \gamma(q, a) \):

\[
\begin{align*}
s_q & \leftarrow \exists \text{succ}_c^-(s_q \land h_i \land c_{i,i}^{(1)}) \\
c_{i,k}^{(1)} & \leftarrow \exists \text{succ}_c^-(s_q \land h_i \land c_{i,j}^{(1)}) \\
h_{i,D} & \leftarrow \exists \text{succ}_r^-(s_q \land h_i \land c_{i,j}^{(1)}).
\end{align*}
\]

Furthermore, the tape cells that are not under the read-write head have their content preserved. Thus, for each \( 1 \leq i < j \leq n \), add

\[
c_{i,j}^{(1)} \leftarrow \exists \text{succ}^- (c_{i,j}^{(1)} \land h_j).
\]

We now identify subtrees that represent accepting computations. For all \( q \in Q_\exists \) and all \( q' \in Q_\forall \) we add the following:

\[
\begin{align*}
s_{\text{acc}} & \leftarrow s_{q_{\text{acc}}} \\
s_{\text{acc}} & \leftarrow s_q \land \exists \text{succ}_c s_{\text{acc}} \land \exists \text{succ}_c s_{\text{acc}}.
\end{align*}
\]

This concludes the reduction. It is not too difficult to check that \( M \) accepts \( w \) iff \( K \) validates \( (C, \{s(a)\}) \).

\section{Conclusions and Outlook}

We have considered the validation of SHACL constraints in the presence of OWL 2 QL ontologies. We have defined the semantics over a carefully constructed notion of canonical model (aka chase) that minimizes the number of fresh successors introduced to satisfy the ontology axioms at each chase step, and we have argued that this semantics is natural and intuitive. We proposed a rewriting algorithm for recursive SHACL constraints with \textit{stratified} negation. It takes as an input a SHACL shapes graph and an ontology, and constructs a new SHACL shapes graph (also with stratified negation) that can be used for sound and complete validation over the data graph alone, without needing to reason about the ontology at validation time. We showed that, under our semantics, validation in the presence of OWL 2 QL ontologies is complete for \textsc{Exptime}, but it remains \textsc{Ptime} complete in data complexity, and hence it is not harder than validation of stratified SHACL alone, without the ontology.

There are several directions for future work. The first is to extend our approach to support more syntactic features of SHACL, like complex path expressions and counting (number restrictions on paths). We believe the mentioned features can be incorporated and supported by our rewriting approach in principle, but it requires a substantial extension. Another direction is to support ontology languages that go beyond OWL 2 QL. We believe our approach can be elegantly generalized to ontologies expressed in \textit{Horn-SHQL}, but it is more challenging to support non-Horn ontology languages. An implementation of our approach also remains for future work. The rewriting algorithm in this paper was meant to demonstrate the principle feasibility of the approach. As presented, our rewriting is best-case exponential since rule 3 in Definition 5.3 forces us to add exponentially many pairs (\( W, H \)). A way to avoid this problem will be needed in order to achieve an efficient implementation of the rewriting. Extending the SHACL fragment to consider \textit{unstratified} negation is also an interesting direction for future work. For this, the \textit{well-founded semantics} for SHACL seems like a promising starting point [2, 6].

\section*{Acknowledgements}

\begin{itemize}
  \item The project leading to this application has received funding from the European Union’s Horizon 2020 research and innovation programme under grant agreement No 101034440.
  \item This work was partially supported by the Wallenberg AI, Autonomous Systems and Software Program (WASP) funded by the Knut and Alice Wallenberg Foundation. In addition, Šimkus was partially supported by the Austrian Science Fund (FWF) project P30873 and Ahmetaj was supported by the FWF and netidee SCIENCE project T1349-N.
\end{itemize}

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